

# Today

- **Memory Layout**
- **Buffer Overflow**
  - Vulnerability
  - Protection
- **Unions**

# x86-64 Linux Memory Layout

*not drawn to scale*

## ■ Stack

- Runtime stack (8MB limit)
- E. g., local variables

## ■ Heap

- Dynamically allocated as needed
- When call `malloc()`, `calloc()`, `new()`

## ■ Data

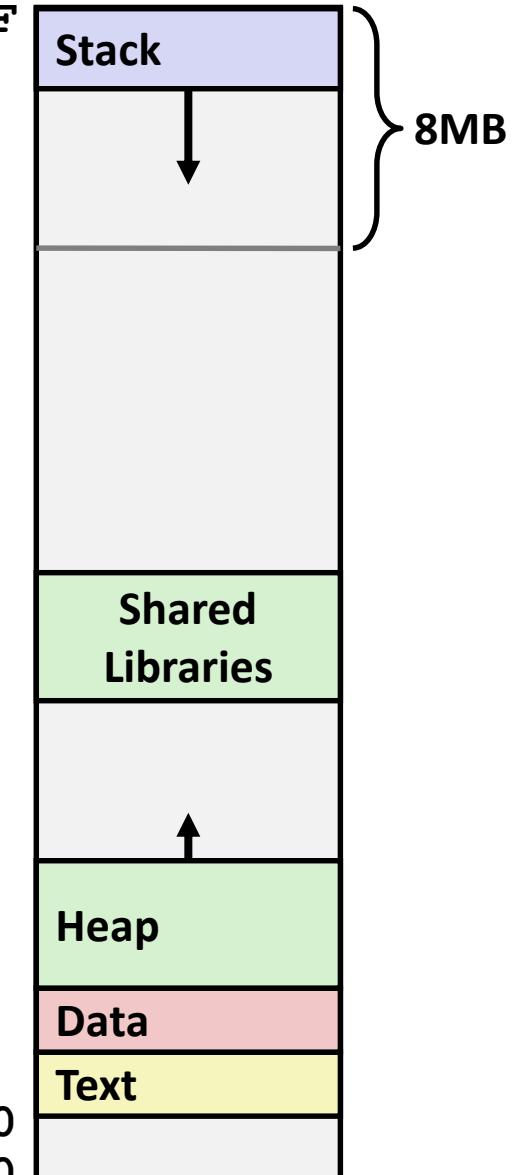
- Statically allocated data
- E.g., global vars, `static` vars, string constants

## ■ Text / Shared Libraries

- Executable machine instructions
- Read-only

Hex Address

400000  
000000



*not drawn to scale*

# Memory Allocation Example

```

char big_array[1L<<24]; /* 16 MB */
char huge_array[1L<<31]; /* 2 GB */

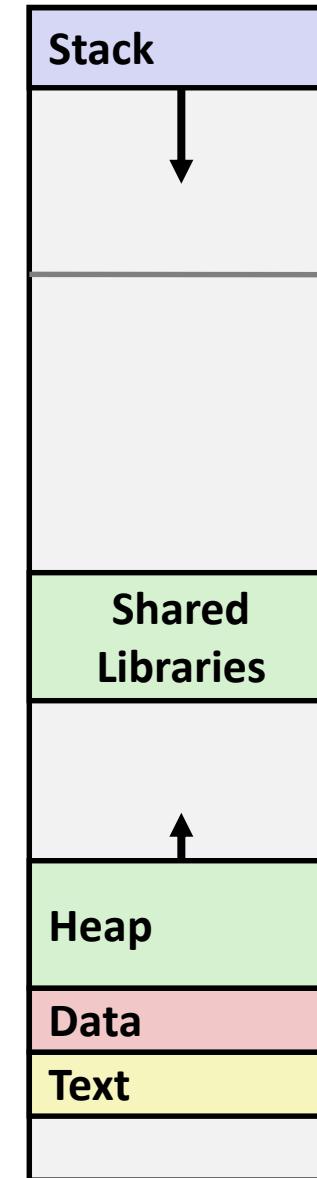
int global = 0;

int useless() { return 0; }

int main (int argc, char** argv)
{
    void *p1, *p2, *p3, *p4;
    int local = 0;
    p1 = malloc(1L << 28); /* 256 MB */
    p2 = malloc(1L << 8); /* 256 B */
    p3 = malloc(1L << 32); /* 4 GB */
    p4 = malloc(1L << 8); /* 256 B */
    /* Some print statements ... */
    return 0;
}

```

*Where does everything go?*

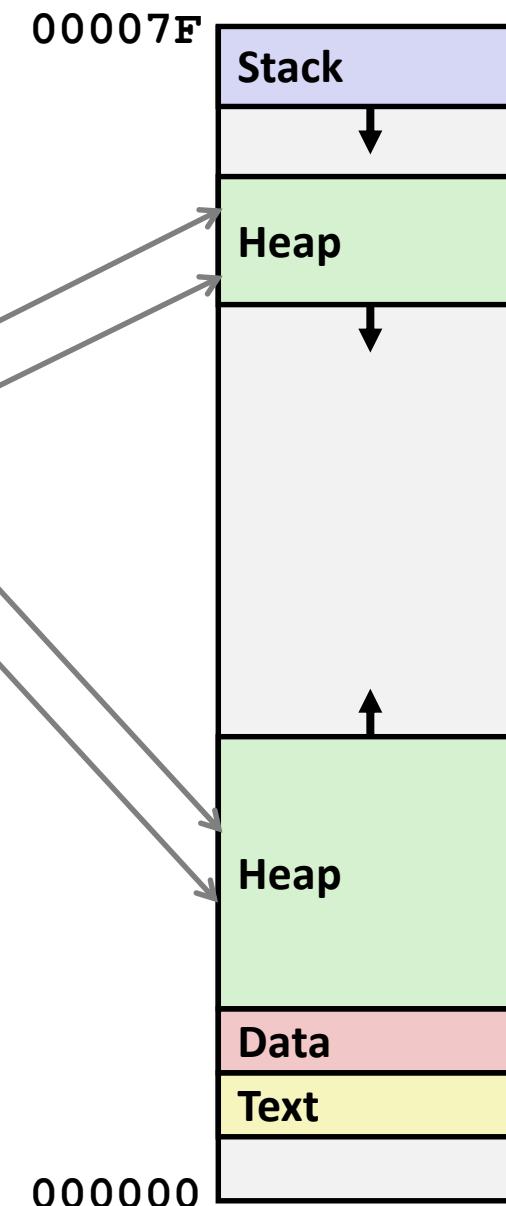


# x86-64 Example Addresses

*address range  $\sim 2^{47}$*

local	0x00007ffe4d3be87c
p1	0x00007f7262a1e010
p3	0x00007f7162a1d010
p4	0x000000008359d120
p2	0x000000008359d010
big_array	0x0000000080601060
huge_array	0x0000000000601060
main()	0x000000000040060c
useless()	0x0000000000400590

*not drawn to scale*



# Today

- Memory Layout
- Buffer Overflow
  - Vulnerability
  - Protection
- Unions

# Recall: Memory Referencing Bug Example

```
typedef struct {
    int a[2];
    double d;
} struct_t;

double fun(int i) {
    volatile struct_t s;
    s.d = 3.14;
    s.a[i] = 1073741824; /* Possibly out of bounds */
    return s.d;
}
```

```
fun(0)  ->  3.14
fun(1)  ->  3.14
fun(2)  ->  3.1399998664856
fun(3)  ->  2.00000061035156
fun(4)  ->  3.14
fun(6)  ->  Segmentation fault
```

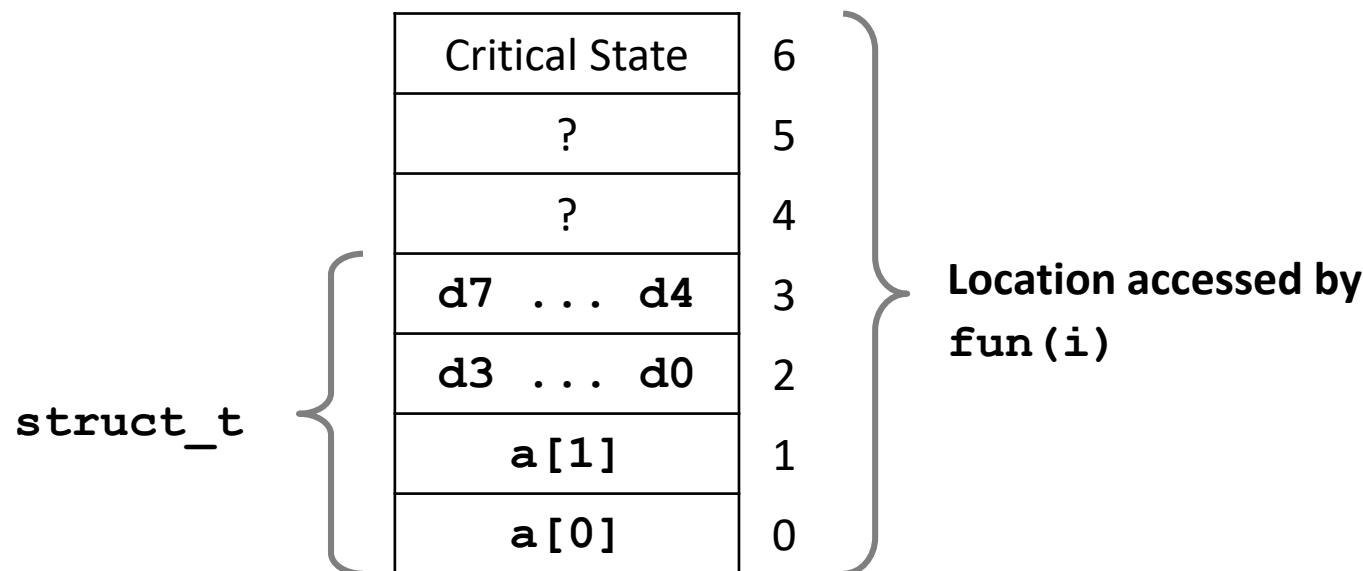
- Result is system specific

# Memory Referencing Bug Example

```
typedef struct {  
    int a[2];  
    double d;  
} struct_t;
```

fun(0)	->	3.14
fun(1)	->	3.14
fun(2)	->	3.1399998664856
fun(3)	->	2.00000061035156
fun(4)	->	3.14
fun(6)	->	Segmentation fault

## Explanation:



# Such problems are a BIG deal

- **Generally called a “buffer overflow”**
  - when exceeding the memory size allocated for an array
- **Why a big deal?**
  - It’s the #1 technical cause of security vulnerabilities
    - #1 overall cause is social engineering / user ignorance
- **Most common form**
  - Unchecked lengths on string inputs
  - Particularly for bounded character arrays on the stack
    - sometimes referred to as stack smashing

# String Library Code

## ■ Implementation of Unix function `gets()`

```
/* Get string from stdin */
char *gets(char *dest)
{
    int c = getchar();
    char *p = dest;
    while (c != EOF && c != '\n') {
        *p++ = c;
        c = getchar();
    }
    *p = '\0';
    return dest;
}
```

- No way to specify limit on number of characters to read
- **Similar problems with other library functions**
  - **strcpy, strcat**: Copy strings of arbitrary length
  - **scanf, fscanf, sscanf**, when given %s conversion specification

# Vulnerable Buffer Code

```
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

←btw, how big  
is big enough?

```
void call_echo() {
    echo();
}
```

```
unix> ./bufdemo-nsp
Type a string: 012345678901234567890123
012345678901234567890123
```

```
unix> ./bufdemo-nsp
Type a string: 0123456789012345678901234
Segmentation Fault
```

# Buffer Overflow Disassembly

echo:

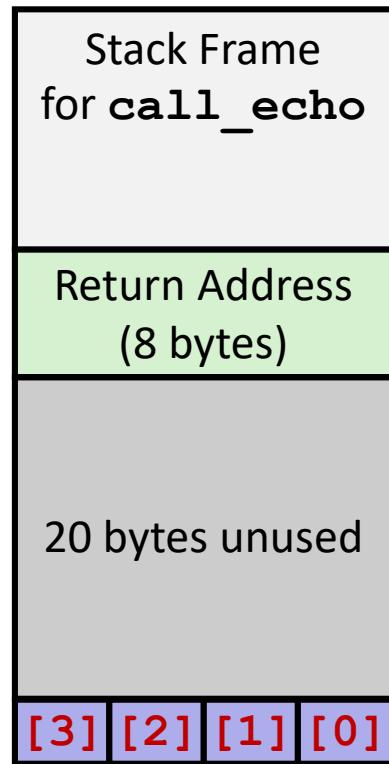
```
00000000004006cf <echo>:  
4006cf: 48 83 ec 18          sub    $0x18,%rsp  
4006d3: 48 89 e7          mov    %rsp,%rdi  
4006d6: e8 a5 ff ff ff      callq   400680 <gets>  
4006db: 48 89 e7          mov    %rsp,%rdi  
4006de: e8 3d fe ff ff      callq   400520 <puts@plt>  
4006e3: 48 83 c4 18          add    $0x18,%rsp  
4006e7: c3                  retq
```

call\_echo:

```
4006e8: 48 83 ec 08          sub    $0x8,%rsp  
4006ec: b8 00 00 00 00      mov    $0x0,%eax  
4006f1: e8 d9 ff ff ff      callq   4006cf <echo>  
4006f6: 48 83 c4 08          add    $0x8,%rsp  
4006fa: c3                  retq
```

# Buffer Overflow Stack

*Before call to gets*

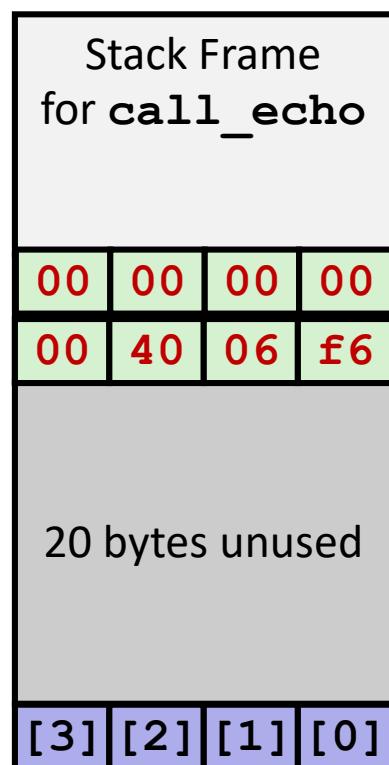


```
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

```
echo:
    subq $24, %rsp
    movq %rsp, %rdi
    call gets
    . . .
```

# Buffer Overflow Stack Example

*Before call to gets*



```
void echo()
{
    char buf[4];
    gets(buf);
    ...
}
```

```
echo:
    subq $24, %rsp
    movq %rsp, %rdi
    call gets
    ...
}
```

`call_echo:`

```
...
4006f1: callq 4006cf <echo>
4006f6: add    $0x8,%rsp
...
```

# Buffer Overflow Stack Example #1

*After call to gets*

Stack Frame for <code>call_echo</code>			
00	00	00	00
00	40	06	f6
00	32	31	30
39	38	37	36
35	34	33	32
31	30	39	38
37	36	35	34
33	32	31	30

```
void echo()
{
    char buf[4];
    gets(buf);
    ...
}
```

```
echo:
    subq $24, %rsp
    movq %rsp, %rdi
    call gets
    ...
}
```

`call_echo:`

```
...
4006f1: callq 4006cf <echo>
4006f6: add    $0x8,%rsp
...
```

`buf ← %rsp`

```
unix>./bufdemo-nsp
Type a string: 01234567890123456789012
01234567890123456789012
```

“01234567890123456789012\0”

**Overflowed buffer, but did not corrupt state**

# Buffer Overflow Stack Example #2

*After call to gets*

Stack Frame for <code>call_echo</code>			
00	00	00	00
00	40	00	34
33	32	31	30
39	38	37	36
35	34	33	32
31	30	39	38
37	36	35	34
33	32	31	30

```
void echo()
{
    char buf[4];
    gets(buf);
    ...
}
```

```
echo:
    subq $24, %rsp
    movq %rsp, %rdi
    call gets
    ...
}
```

`call_echo:`

```
...
4006f1: callq 4006cf <echo>
4006f6: add    $0x8,%rsp
...
```

`buf` ← `%rsp`

```
unix>./bufdemo-nsp
Type a string: 0123456789012345678901234
Segmentation Fault
```

“012345678901234567890123**4\0**”

**Overflowed buffer and corrupted return pointer**

# Buffer Overflow Stack Example #3

*After call to gets*

Stack Frame for <code>call_echo</code>			
00	00	00	00
00	40	06	00
33	32	31	30
39	38	37	36
35	34	33	32
31	30	39	38
37	36	35	34
33	32	31	30

```
void echo()
{
    char buf[4];
    gets(buf);
    ...
}
```

```
echo:
    subq $24, %rsp
    movq %rsp, %rdi
    call gets
    ...
}
```

`call_echo:`

```
...
4006f1: callq 4006cf <echo>
4006f6: add    $0x8,%rsp
...
```

`buf` ← `%rsp`

```
unix>./bufdemo-nsp
Type a string: 012345678901234567890123
012345678901234567890123
```

“012345678901234567890123\0”

**Overflowed buffer, corrupted return pointer, but program seems to work!**

# Buffer Overflow Stack Example #3 Explained

*After call to gets*

Stack Frame for call echo			
00	00	00	00
00	48	83	80
00	00	00	00
00	40	06	00
33	32	31	30
39	38	37	36
35	34	33	32
31	30	39	38
37	36	35	34
33	32	31	30

buf ← %rsp

**register\_tm\_clones:**

```
...
400600:    mov    %rsp,%rbp
400603:    mov    %rax,%rdx
400606:    shr    $0x3f,%rdx
40060a:    add    %rdx,%rax
40060d:    sar    %rax
400610:    jne    400614
400612:    pop    %rbp
400613:    retq
```

```
void call_echo() {
    echo();
}
```

“Returns” to unrelated code

Lots of things happen, without modifying critical state

Eventually executes `retq` back to `main` – `call_echo` has no local variables

# Stack Smashing Attacks

```

void P() {
    Q();
    ...
}

```

return address  
**A**

```

int Q() {
    char buf[64];
    gets(buf);
    ...
    return ...;
}

```

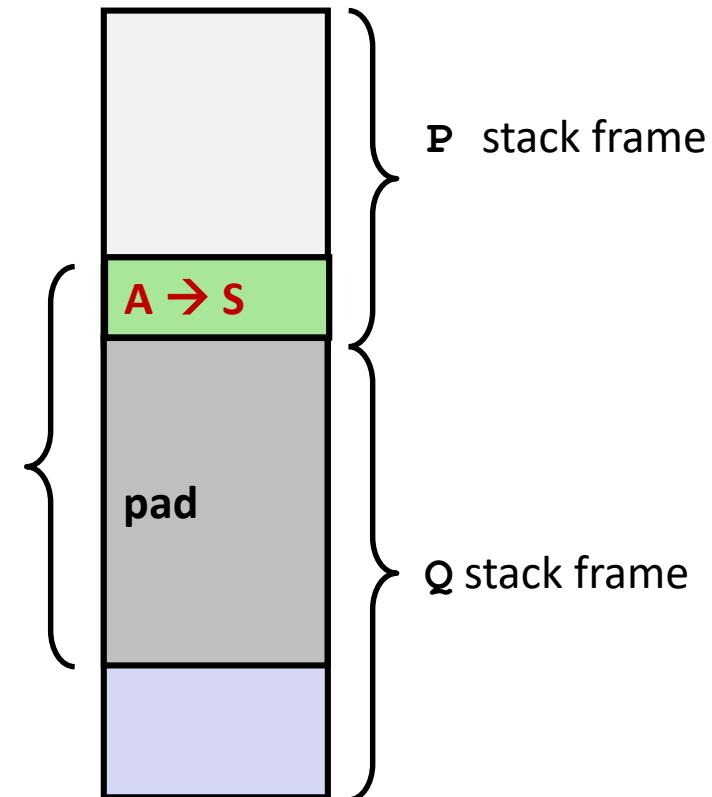
```

void S() {
    /* Something
       unexpected */
    ...
}

```

data written by `gets()`

Stack after call to `gets()`



- Overwrite normal return address A with address of some other code S
- When Q executes `ret`, will jump to other code

# Crafting Smashing String

Stack Frame for call echo			
00	00	00	00
00	48	83	80
00	00	00	00
00	40	08	83

```
int echo() {
    char buf[4];
    gets(buf);
    ...
    return ...;
}
```

← %rsp

24 bytes

*Target Code*

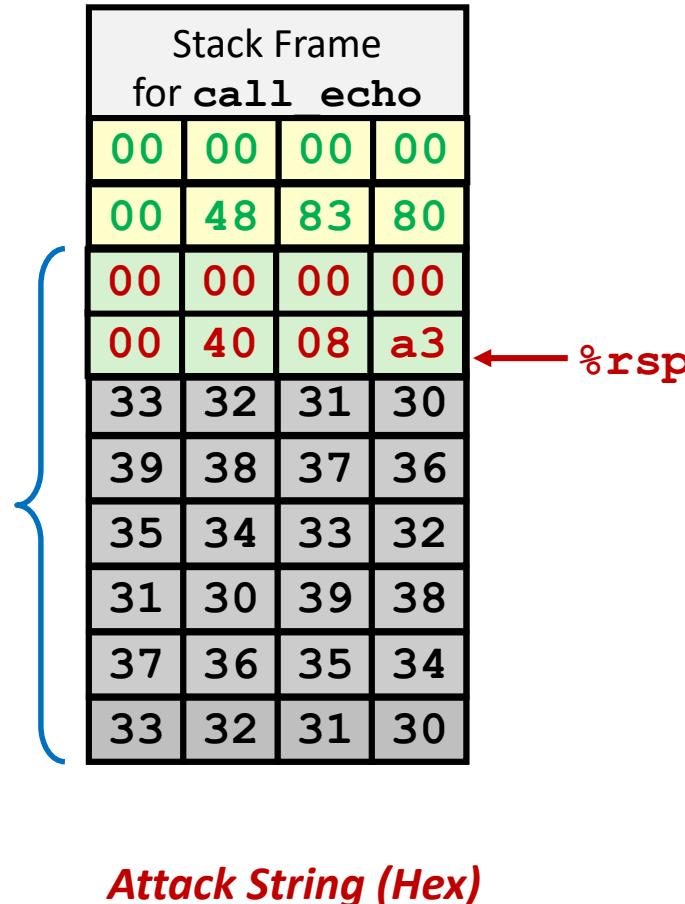
```
void smash() {
    printf("I've been smashed!\n");
    exit(0);
}
```

00000000004008a3 <smash>:  
4008a3: 48 83 ec 08

*Attack String (Hex)*

30	31	32	33	34	35	36	37	38	39	30	31	32	33	34	35	36	37	38	39	30	31	32	33
a3	08	40	00	00	00	00	00	00	00														

# Smashing String Effect



**Target Code**

```
void smash() {
    printf("I've been smashed!\n");
    exit(0);
}
```

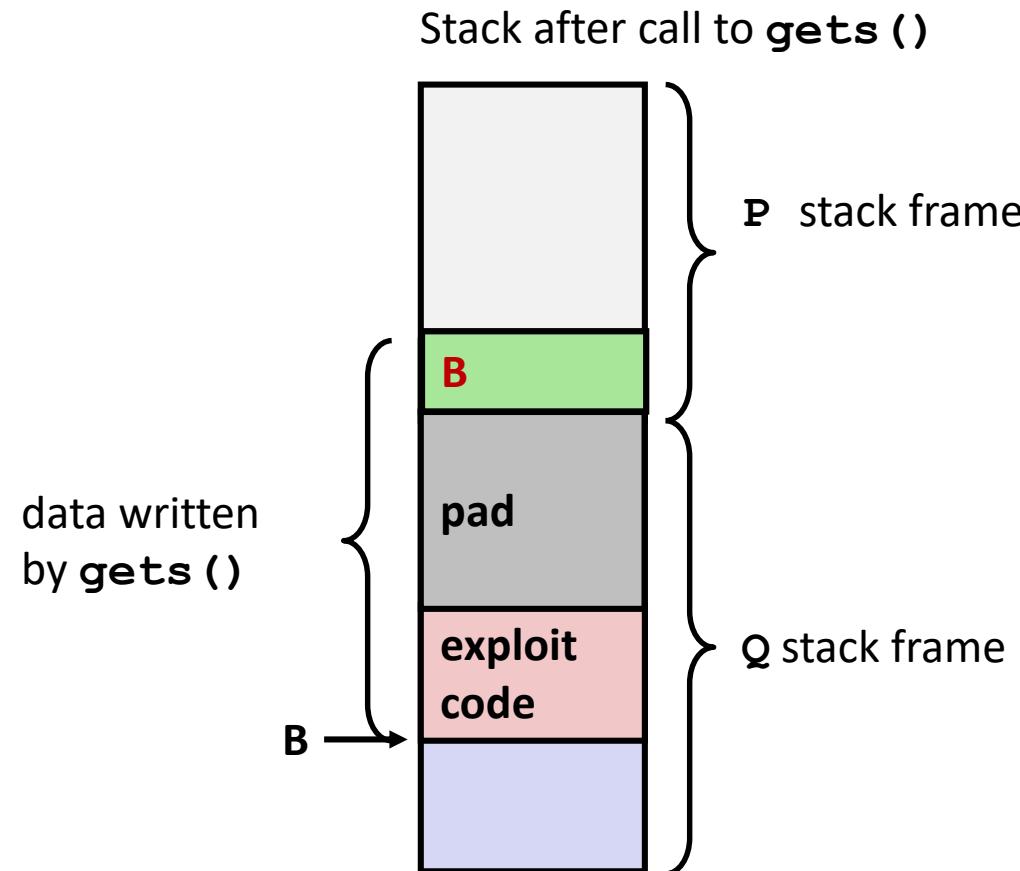
00000000004008a3 <smash>:  
4008a3: 48 83 ec 08

# Code Injection Attacks

```

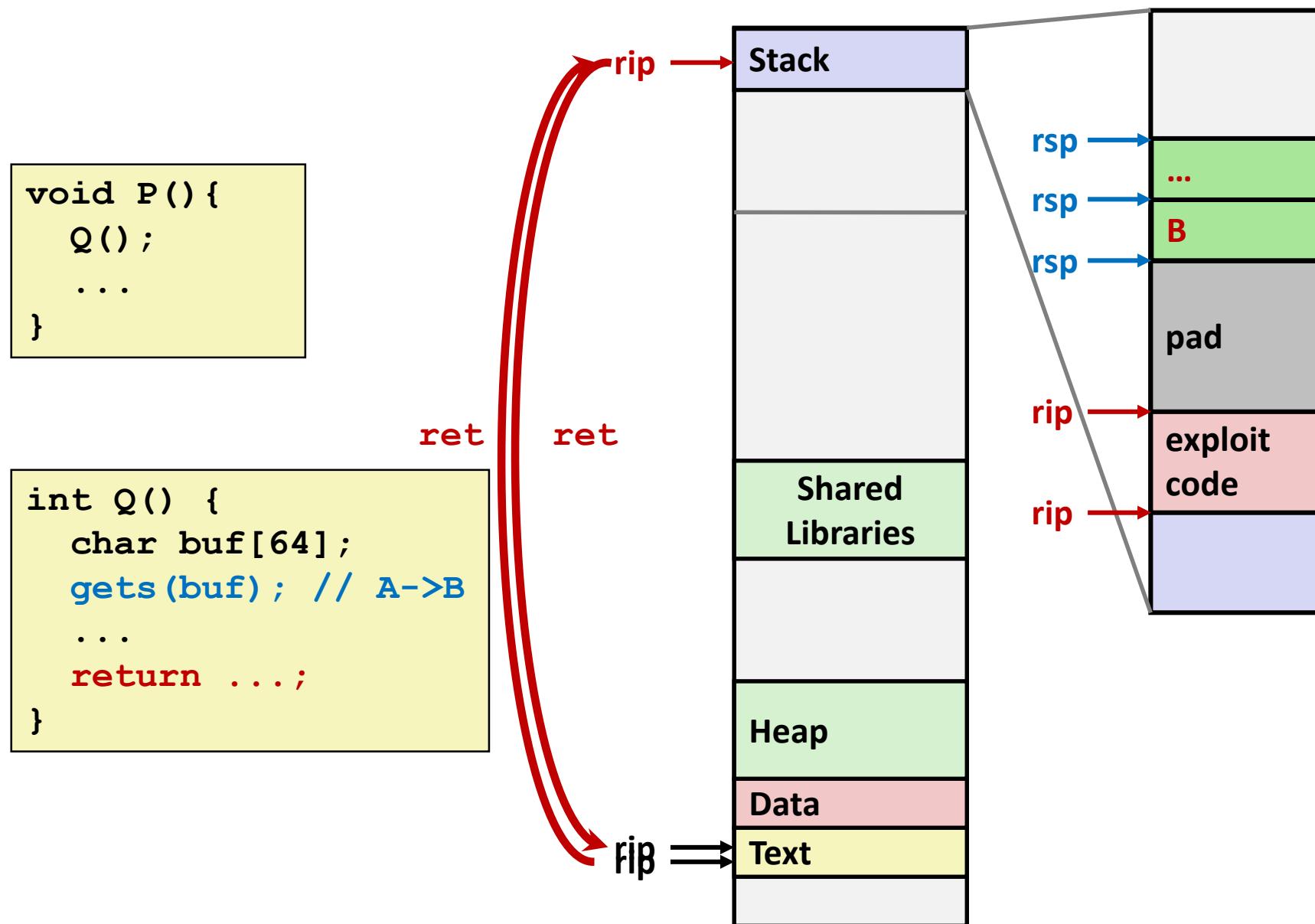
void P() {
    Q();
    ...
}
int Q() {
    char buf[64];
    gets(buf);
    ...
    return ...;
}

```



- Input string contains byte representation of executable code
- Overwrite return address A with address of buffer B
- When Q executes `ret`, will jump to exploit code

# How Does The Attack Code Execute?



# What To Do About Buffer Overflow Attacks

- Avoid overflow vulnerabilities
- Employ system-level protections
- Have compiler use “stack canaries”
- Lets talk about each...

# 1. Avoid Overflow Vulnerabilities in Code (!)

```
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    fgets(buf, 4, stdin);
    puts(buf);
}
```

- For example, use library routines that limit string lengths
  - **fgets** instead of **gets**
  - **strncpy** instead of **strcpy**
  - Don't use **scanf** with **%s** conversion specification
    - Use **fgets** to read the string
    - Or use **%ns** where **n** is a suitable integer

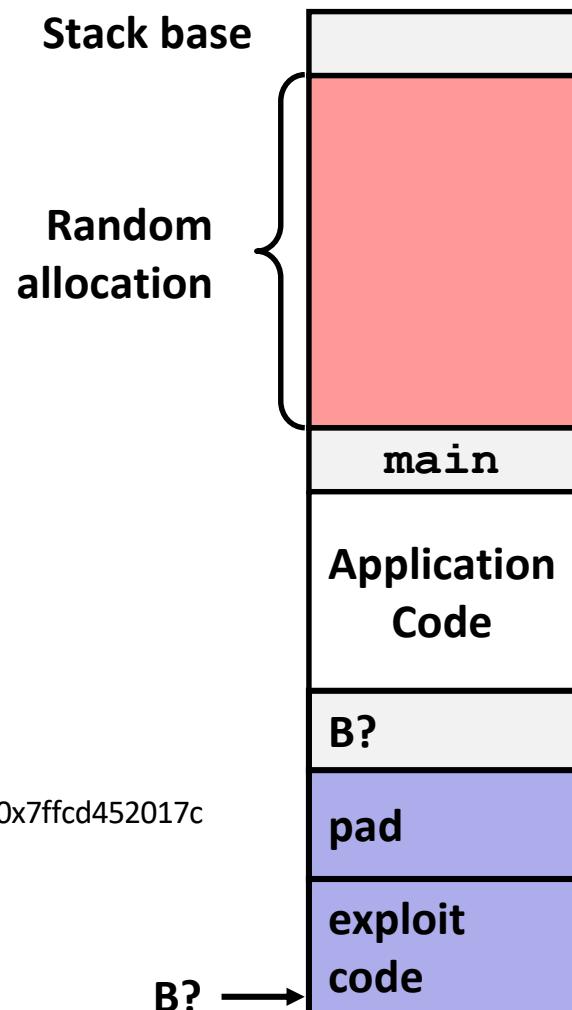
## 2. System-Level Protections can help

### ■ Randomized stack offsets

- At start of program, allocate random amount of space on stack
- Shifts stack addresses for entire program
- Makes it difficult for hacker to predict beginning of inserted code
- E.g.: 5 executions of memory allocation code

local      0x7ffe4d3be87c    0x7fff75a4f9fc    0x7ffeadb7c80c    0x7ffeaea2fdac    0x7ffcd452017c

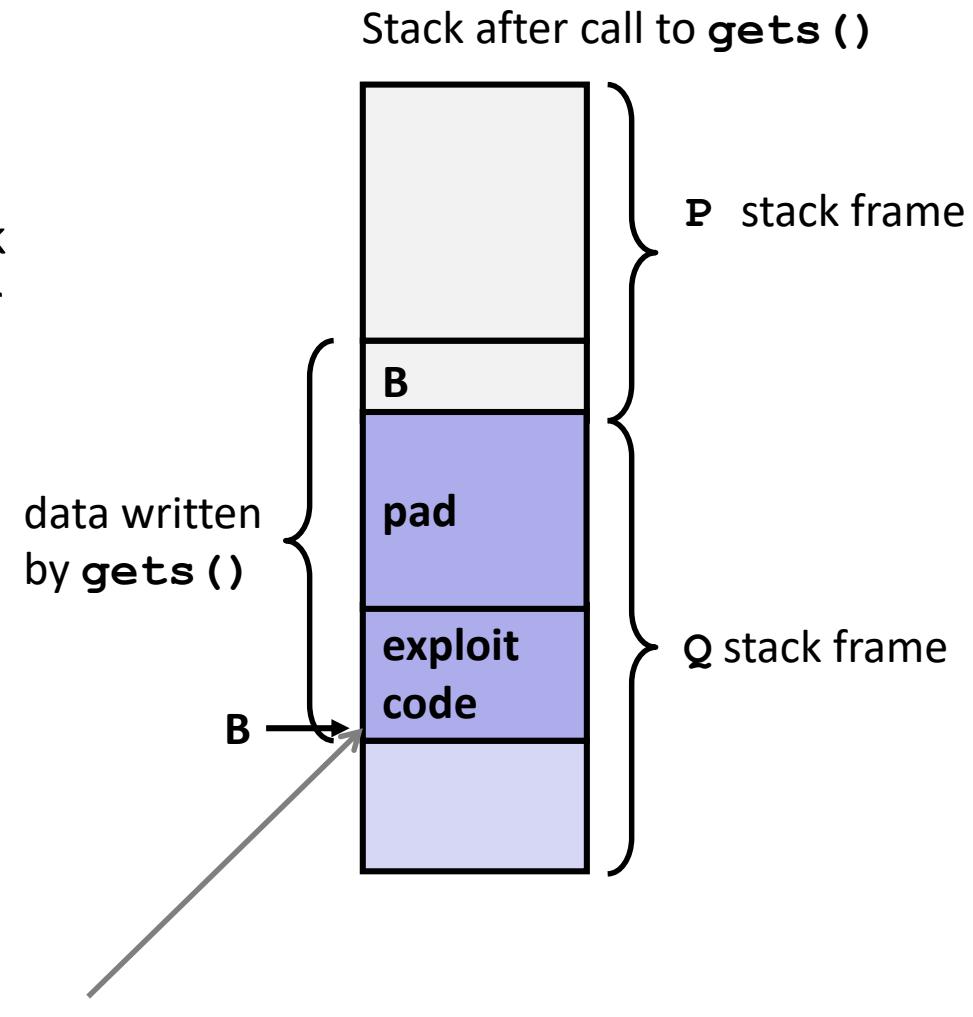
- Stack repositioned each time program executes



## 2. System-Level Protections can help

### ■ Nonexecutable code segments

- In traditional x86, can mark region of memory as either “read-only” or “writeable”
  - Can execute anything readable
- x86-64 added explicit “execute” permission
- Stack marked as non-executable



# 3. Stack Canaries can help

## ■ Idea

- Place special value (“canary”) on stack just beyond buffer
- Check for corruption before exiting function

## ■ GCC Implementation

- **-fstack-protector**
- Now the default (disabled earlier)

```
unix>./bufdemo-sp
Type a string: 0123456
0123456
```

```
unix>./bufdemo-sp
Type a string: 01234567
*** stack smashing detected ***
```

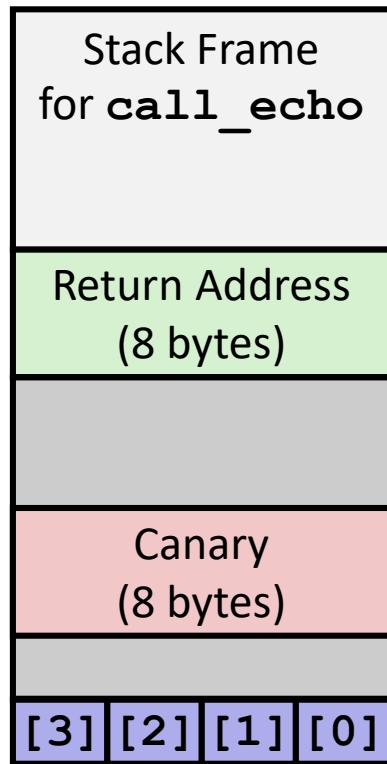
# Protected Buffer Disassembly

echo:

```
40072f: sub    $0x18,%rsp
400733: mov    %fs:0x28,%rax
40073c: mov    %rax,0x8(%rsp)
400741: xor    %eax,%eax
400743: mov    %rsp,%rdi
400746: callq  4006e0 <gets>
40074b: mov    %rsp,%rdi
40074e: callq  400570 <puts@plt>
400753: mov    0x8(%rsp),%rax
400758: xor    %fs:0x28,%rax
400761: je    400768 <echo+0x39>
400763: callq  400580 <__stack_chk_fail@plt>
400768: add    $0x18,%rsp
40076c: retq
```

# Setting Up Canary

*Before call to gets*

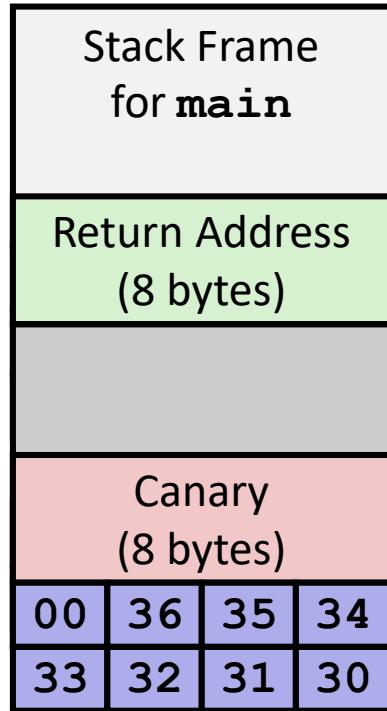


```
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

```
echo:
    . . .
    movq    %fs:40, %rax # Get canary
    movq    %rax, 8(%rsp) # Place on stack
    xorl    %eax, %eax   # Erase canary
    . . .
```

# Checking Canary

*After call to gets*



```
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

**Input: 0123456**

```
echo:
    .
    .
    movq    8(%rsp), %rax      # Retrieve from stack
    xorq    %fs:40, %rax      # Compare to canary
    je     .L6                  # If same, OK
    call   __stack_chk_fail    # FAIL
```