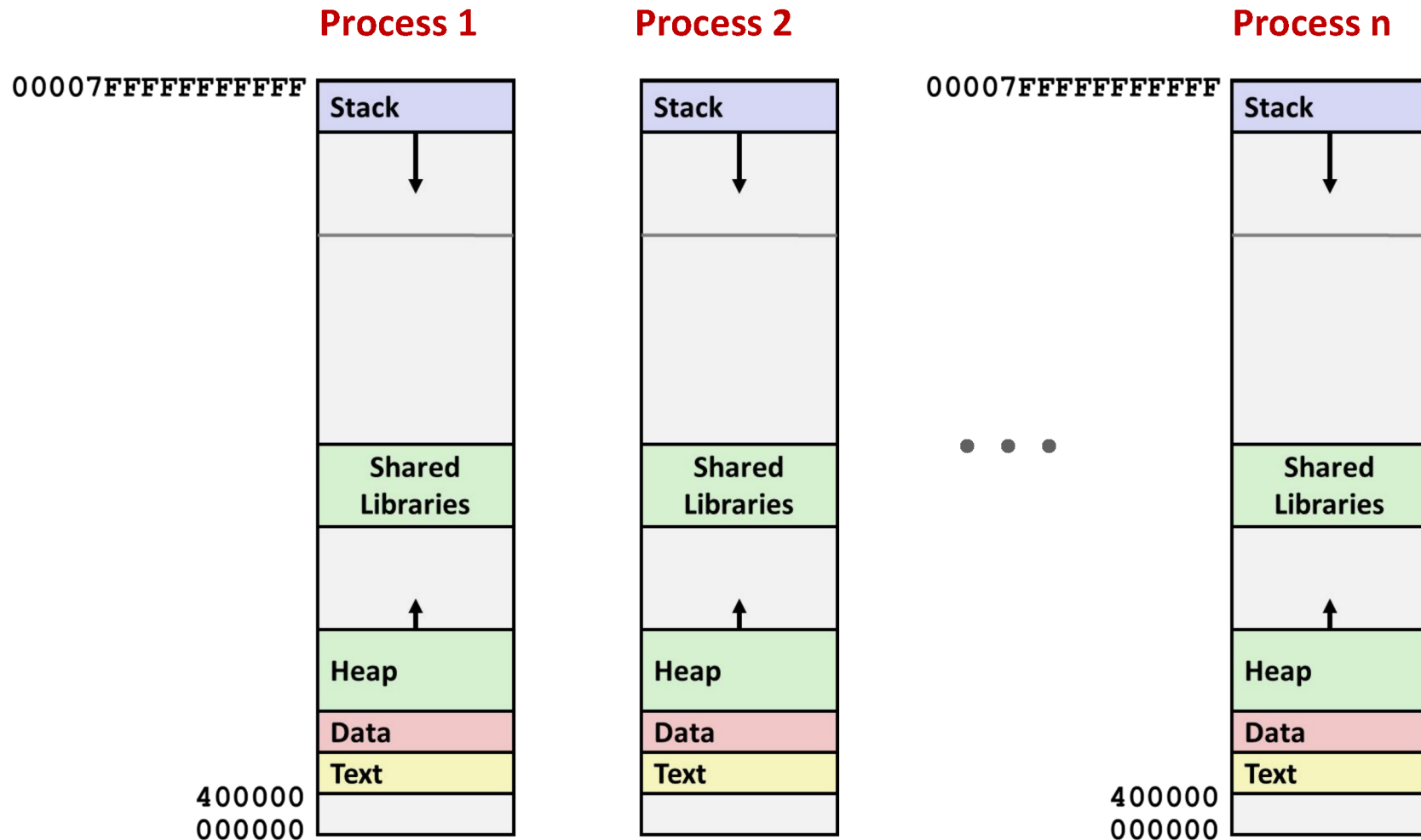


Hmmm, How Does This Work?!

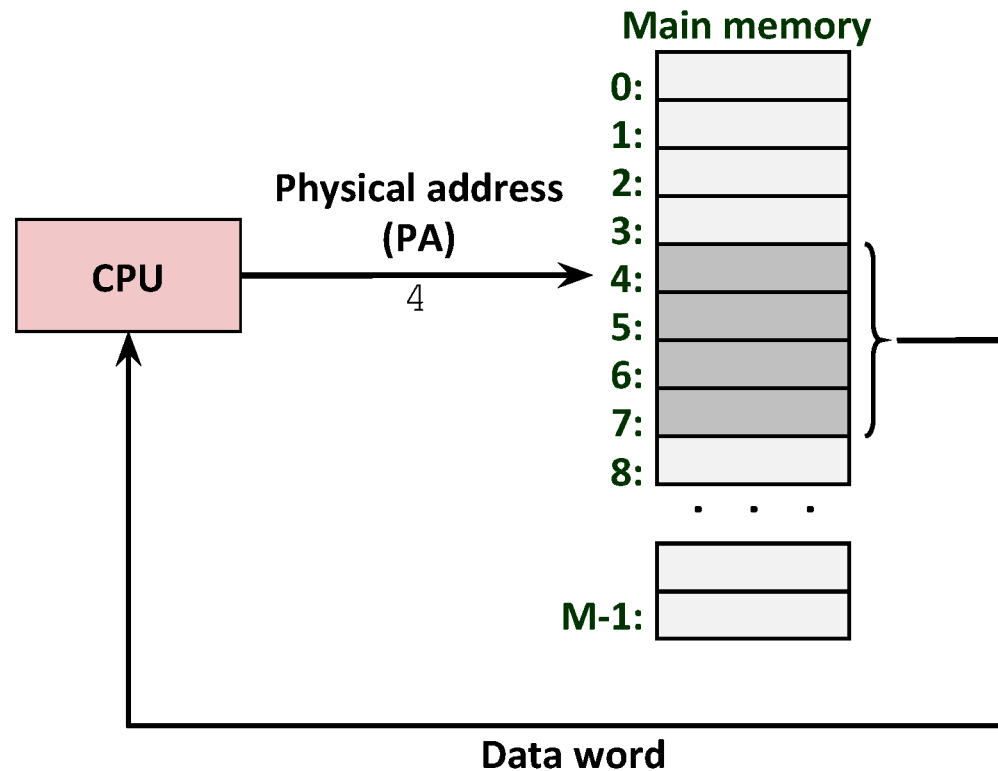


Solution: Virtual Memory (today and next lecture)

Today

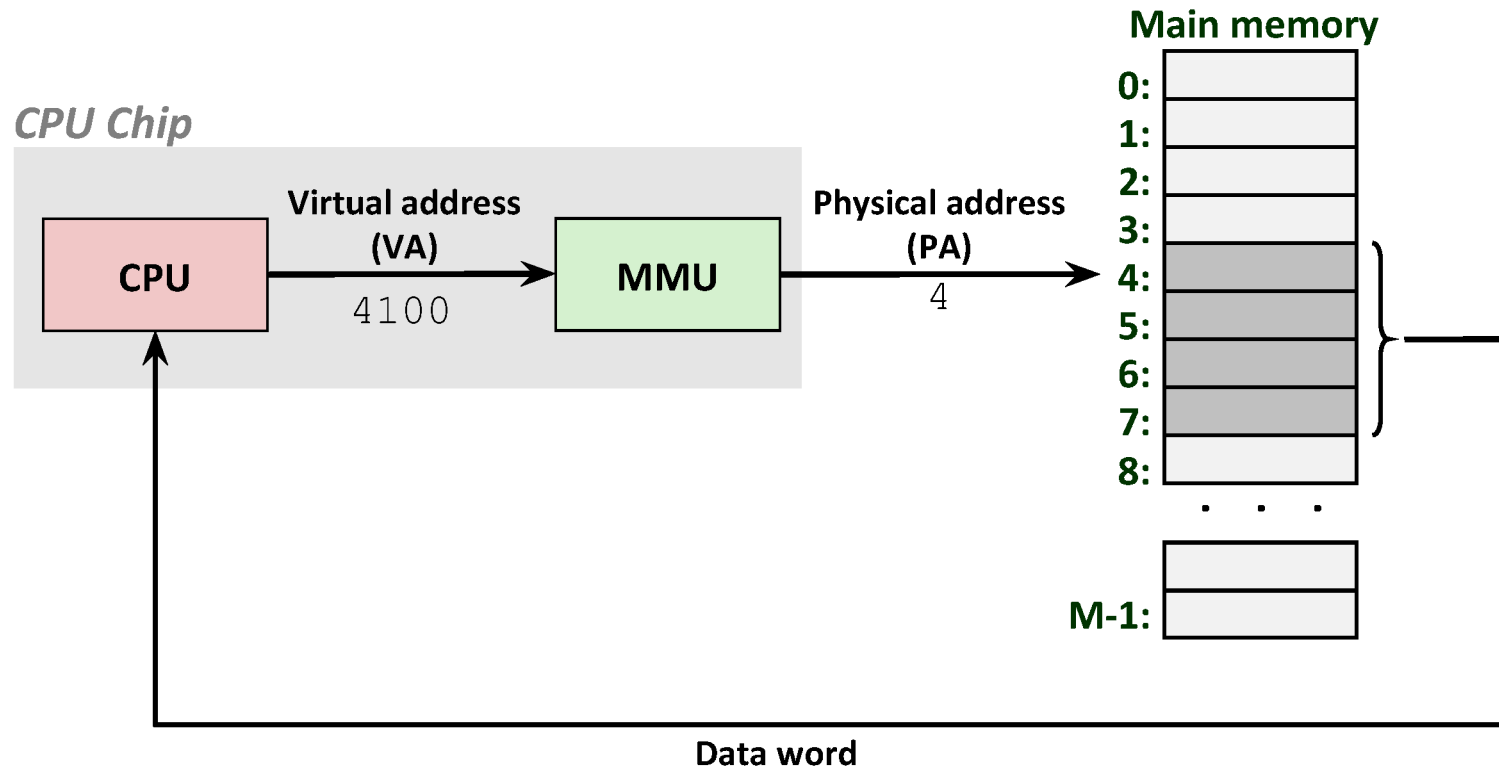
- **Address spaces**
- VM as a tool for caching
- VM as a tool for memory management
- VM as a tool for memory protection
- Address translation

A System Using Physical Addressing



- Used in “simple” systems like embedded microcontrollers in devices like cars, elevators, and digital picture frames

A System Using Virtual Addressing



- Used in all modern servers, laptops, and smart phones
- One of the great ideas in computer science

Address Spaces

- **Linear address space:** Ordered set of contiguous non-negative integer addresses:
 $\{0, 1, 2, 3 \dots \}$
- **Virtual address space:** Set of $N = 2^n$ virtual addresses
 $\{0, 1, 2, 3, \dots, N-1\}$
- **Physical address space:** Set of $M = 2^m$ physical addresses
 $\{0, 1, 2, 3, \dots, M-1\}$

Why Virtual Memory (VM)?

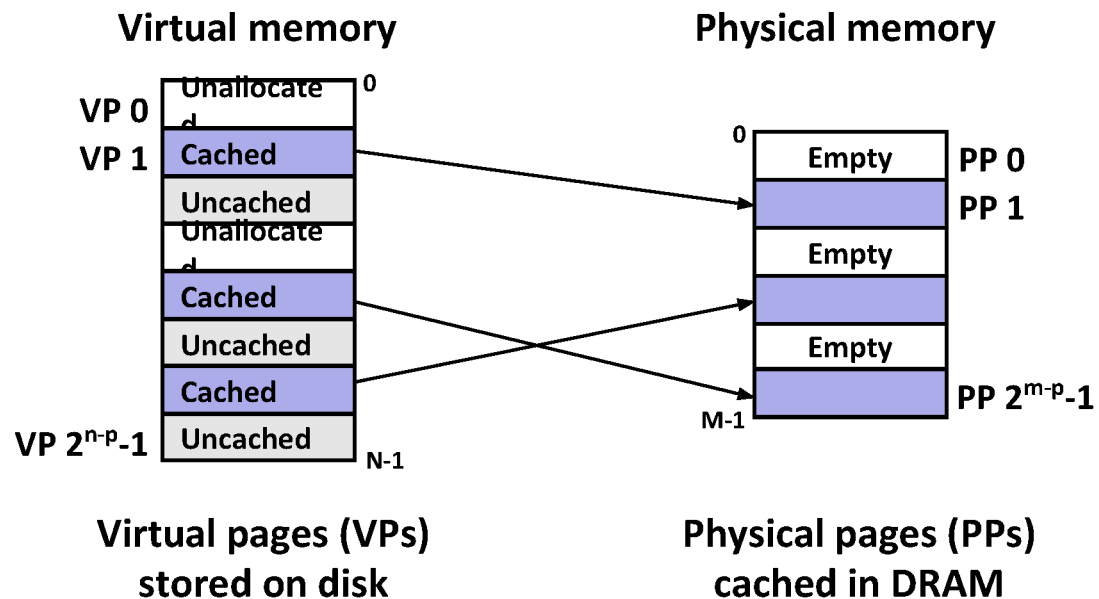
- **Uses main memory efficiently**
 - Use DRAM as a cache for parts of a virtual address space
- **Simplifies memory management**
 - Each process gets the same uniform linear address space
- **Isolates address spaces**
 - One process can't interfere with another's memory
 - User program cannot access privileged kernel information and code

Today

- Address spaces
- **VM as a tool for caching**
- VM as a tool for memory management
- VM as a tool for memory protection
- Address translation

VM as a Tool for Caching

- Conceptually, *virtual memory* is an array of N contiguous bytes stored on disk.
- The contents of the array on disk are cached in *physical memory (DRAM cache)*
 - These cache blocks are called *pages* (size is $P = 2^p$ bytes)



Remember: Set Associative Cache

$E = 2$: Two lines per set

Assume: cache block size 8 bytes

Address :

Block
offset

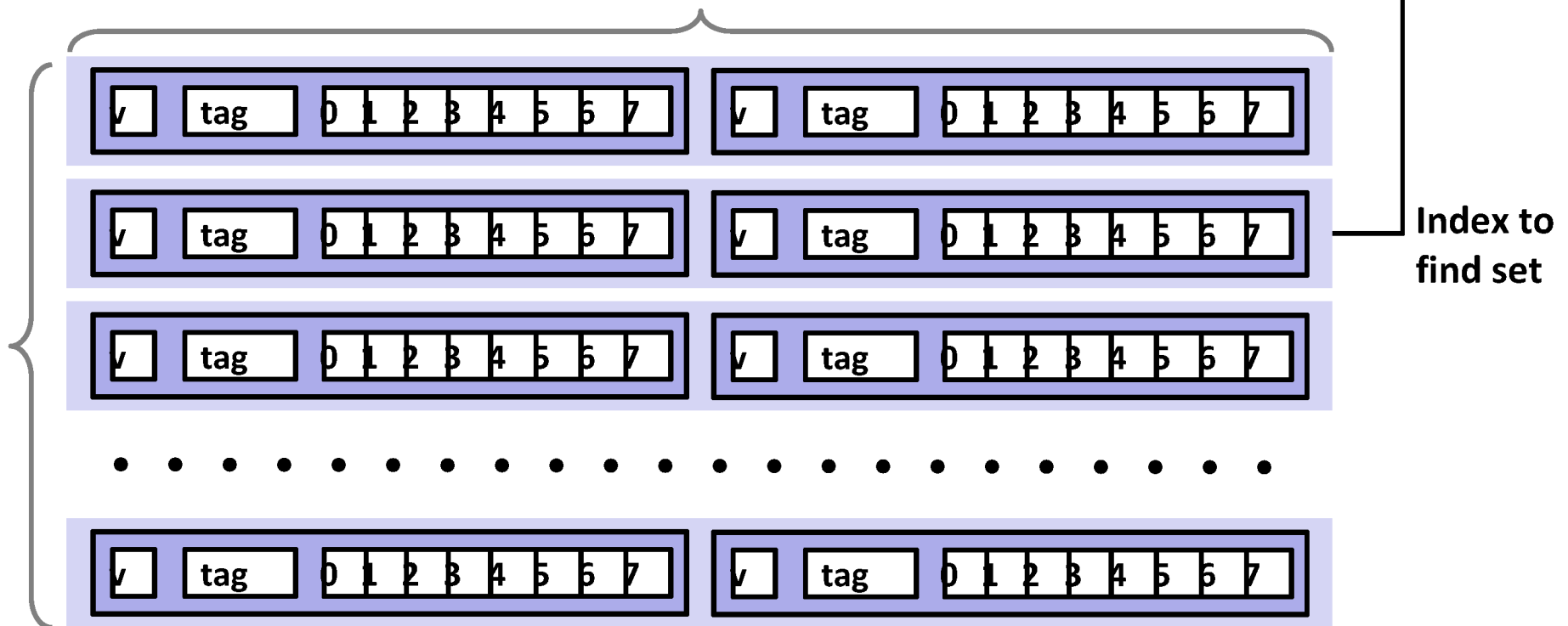


2 lines per set

t bits

0...01

100



S sets

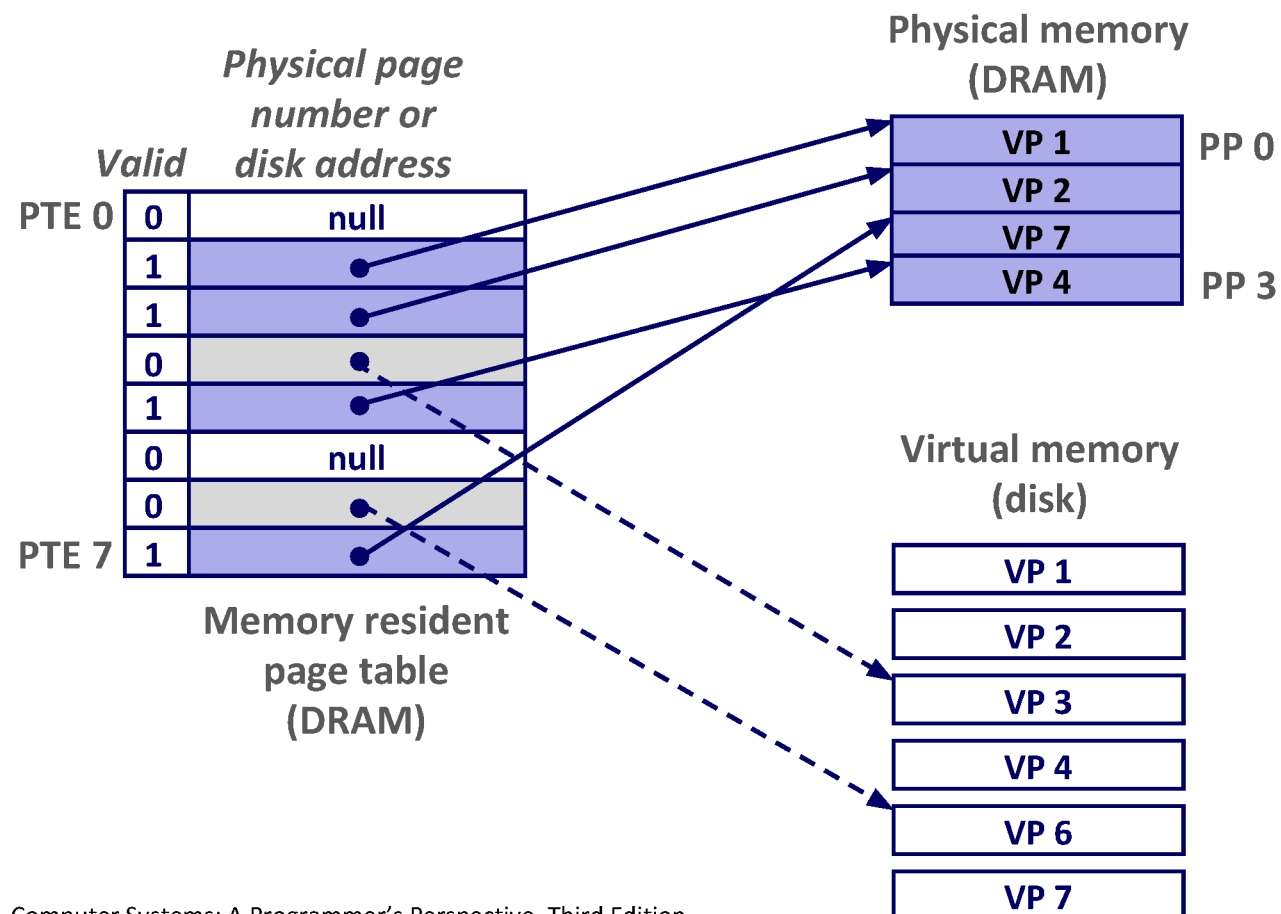
DRAM Cache Organization

- **DRAM cache organization driven by the enormous miss penalty**
 - DRAM is about **10x** slower than SRAM
 - Disk is about **10,000x** slower than DRAM

- **Consequences**
 - Large page (block) size: typically 4 KB, sometimes 4 MB
 - Fully associative
 - Any VP can be placed in any PP
 - Requires a “large” mapping function – different from cache memories
 - Highly sophisticated, expensive replacement algorithms
 - Too complicated and open-ended to be implemented in hardware
 - Write-back rather than write-through

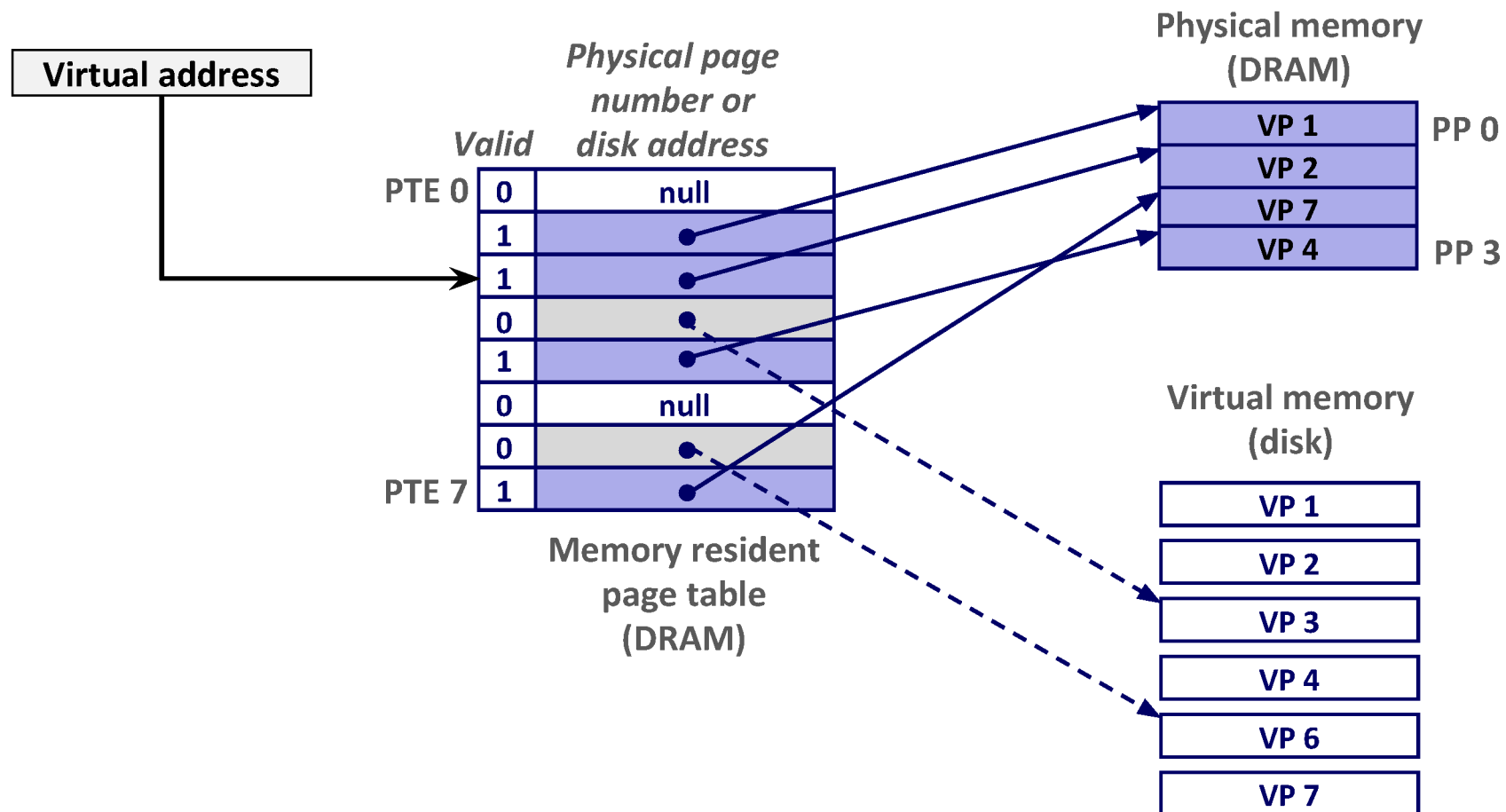
Enabling Data Structure: Page Table

- A *page table* is an array of page table entries (PTEs) that maps virtual pages to physical pages.
 - Per-process kernel data structure in DRAM



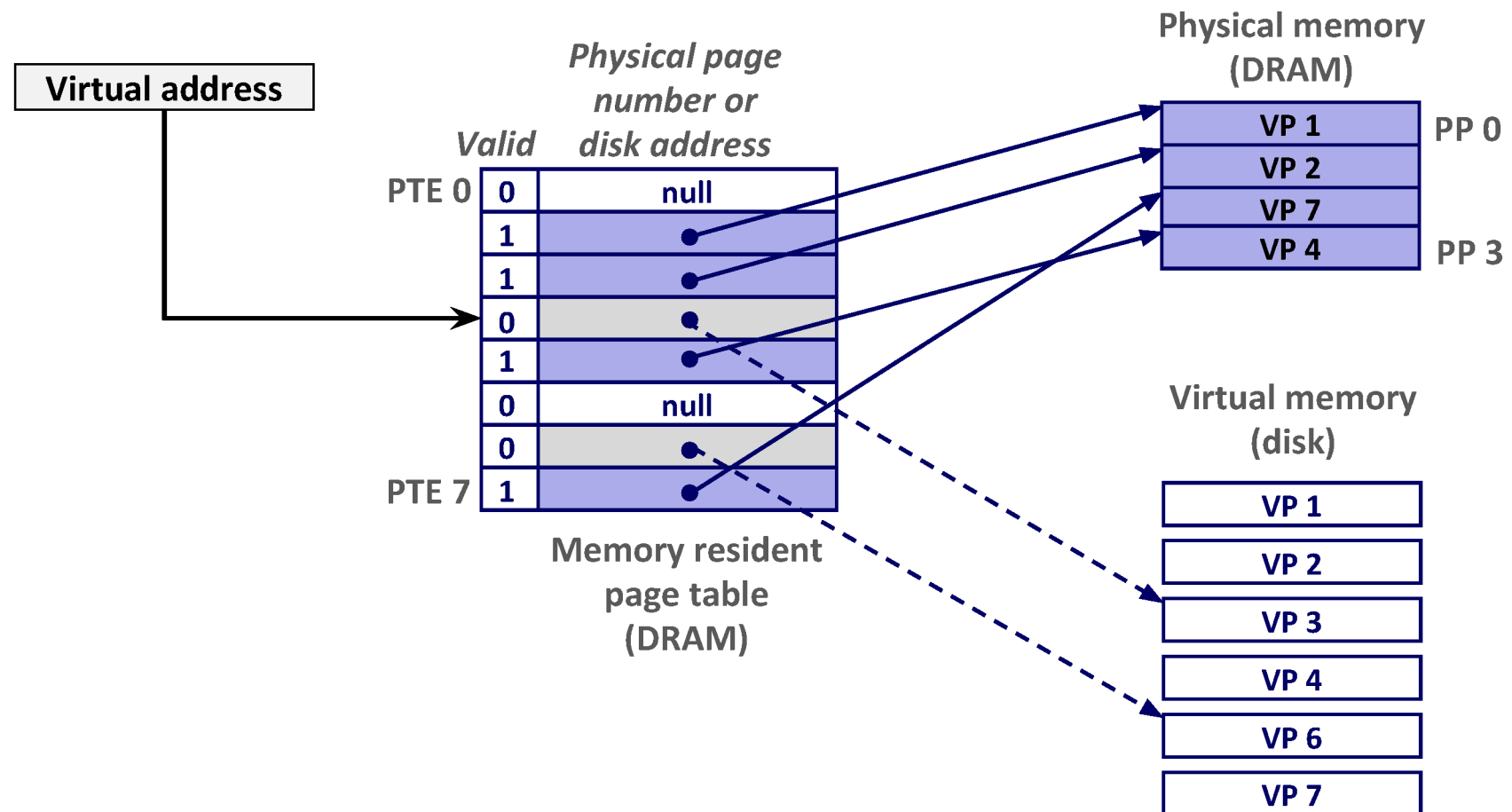
Page Hit

- **Page hit:** reference to VM word that is in physical memory (DRAM cache hit)



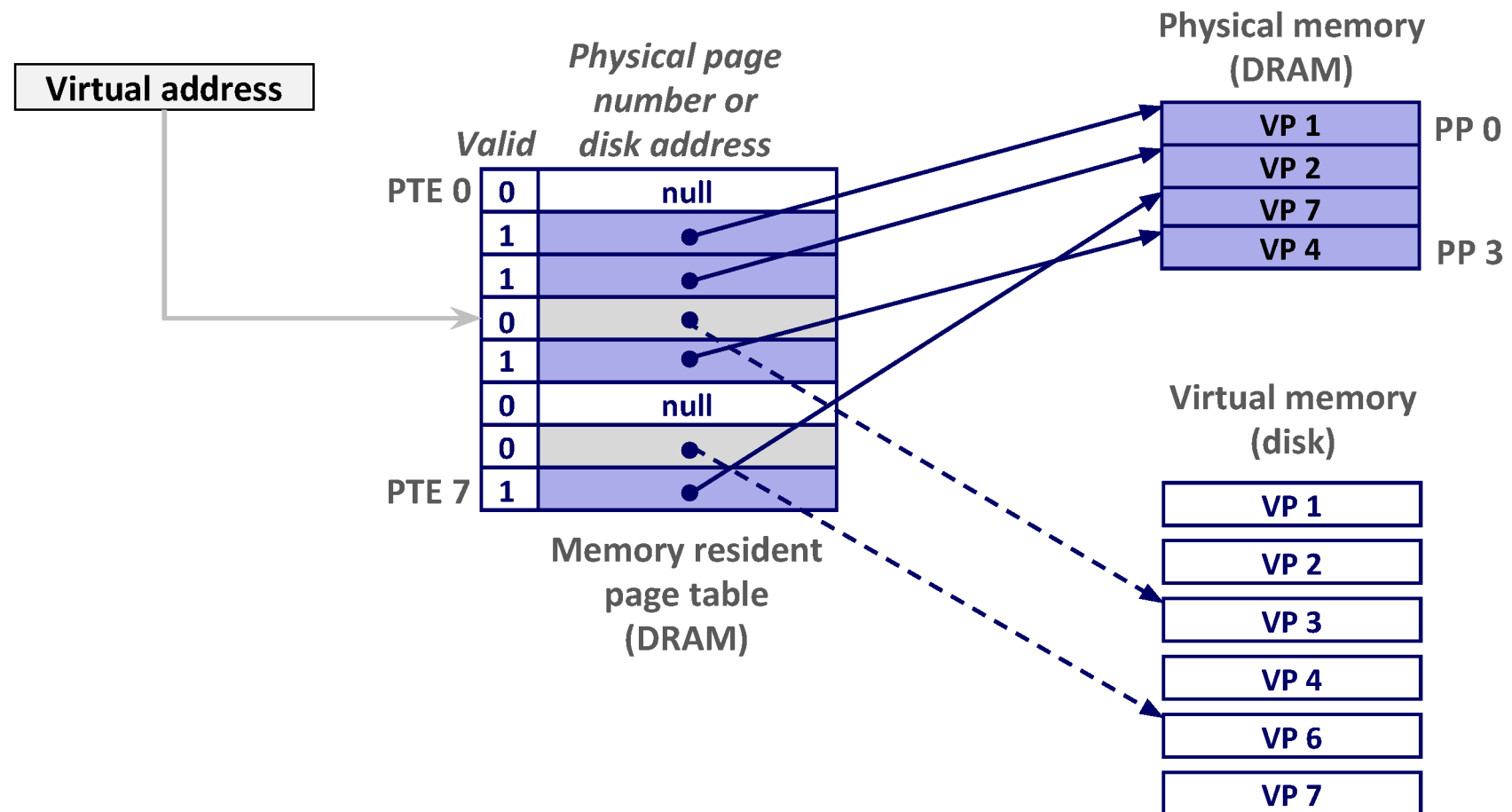
Page Fault

- **Page fault:** reference to VM word that is not in physical memory (DRAM cache miss)



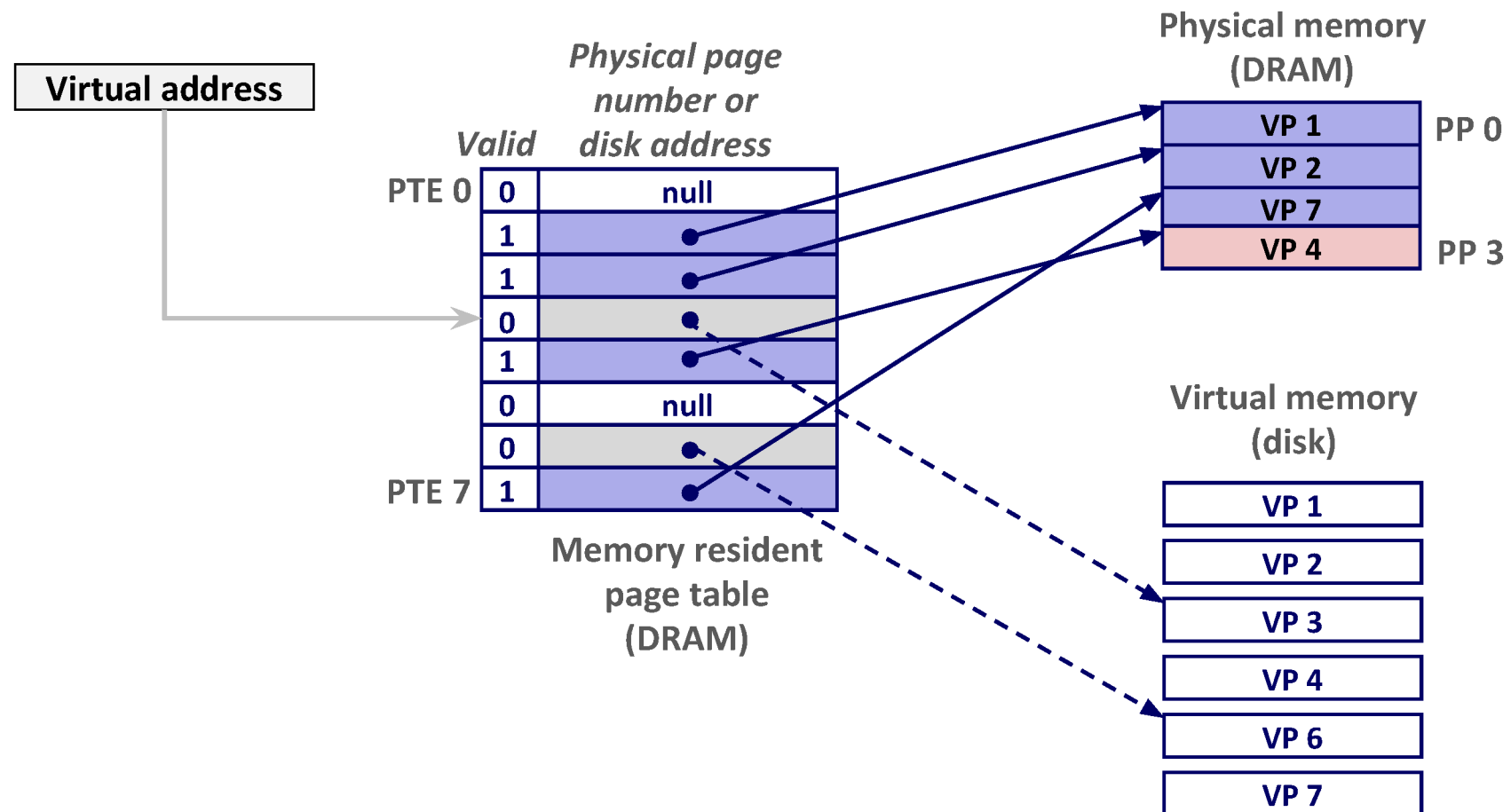
Handling Page Fault

- Page miss causes page fault (an exception)



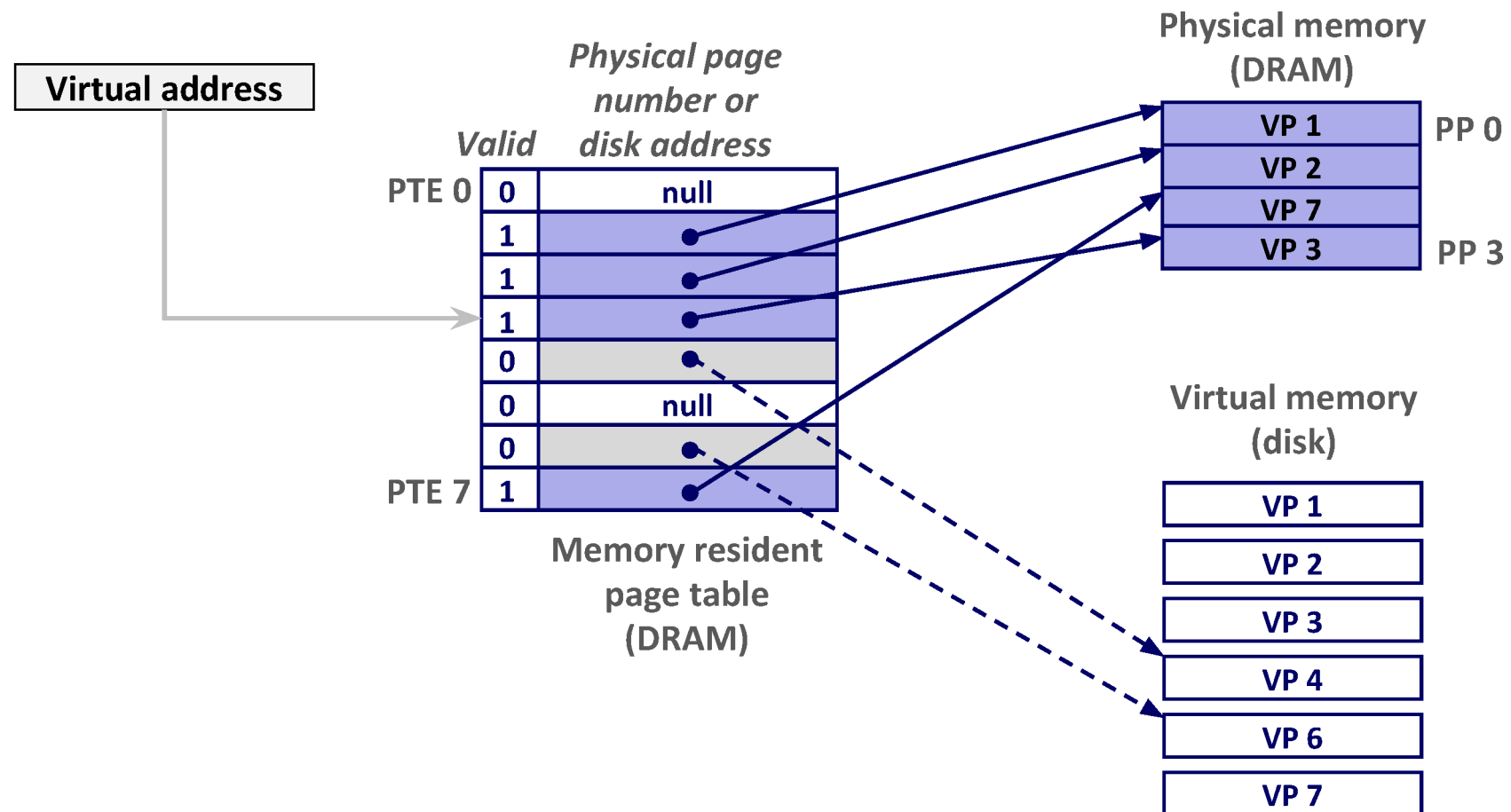
Handling Page Fault

- Page miss causes page fault (an exception)
- Page fault handler selects a victim to be evicted (here VP 4)



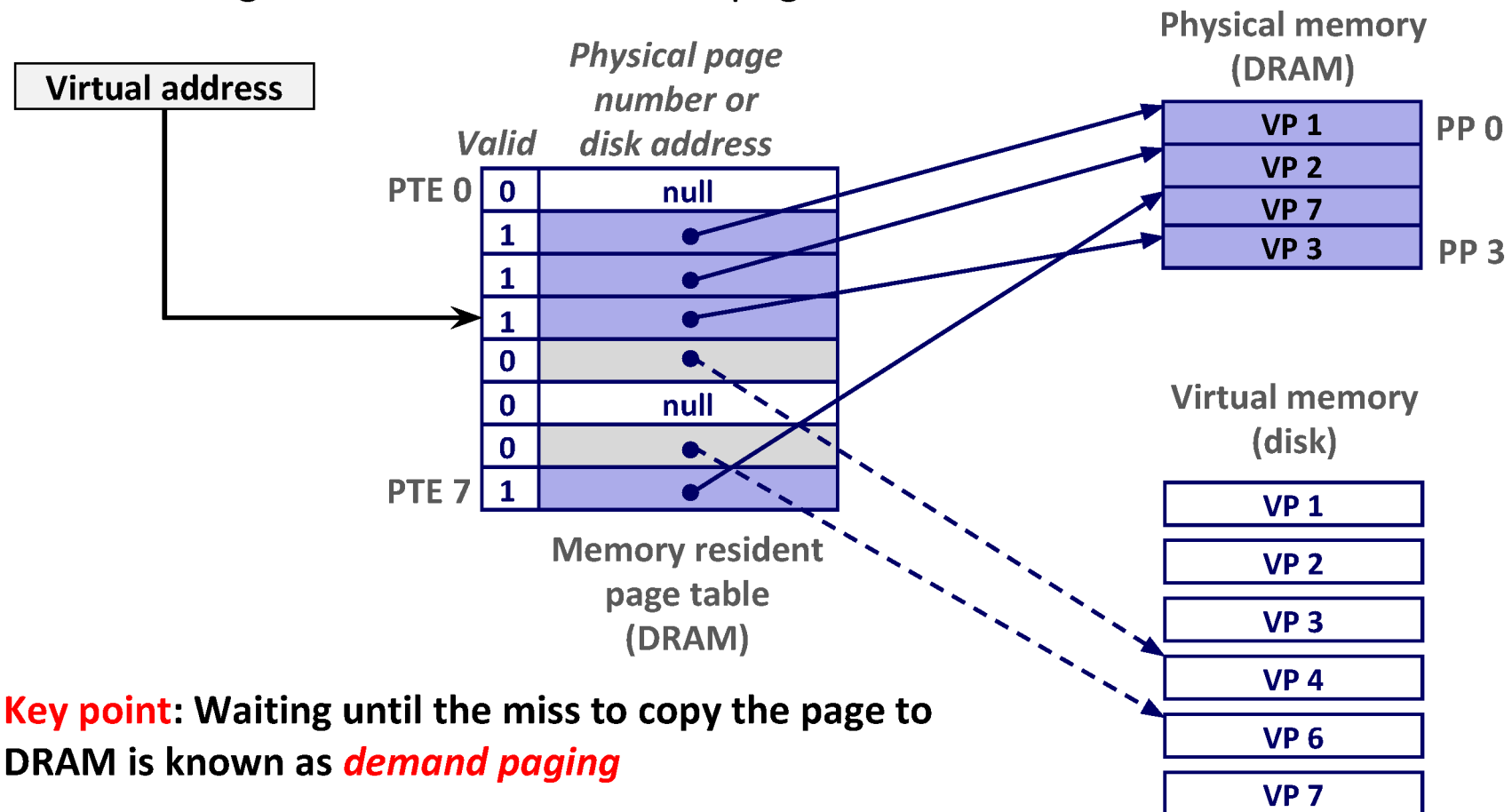
Handling Page Fault

- Page miss causes page fault (an exception)
- Page fault handler selects a victim to be evicted (here VP 4)



Handling Page Fault

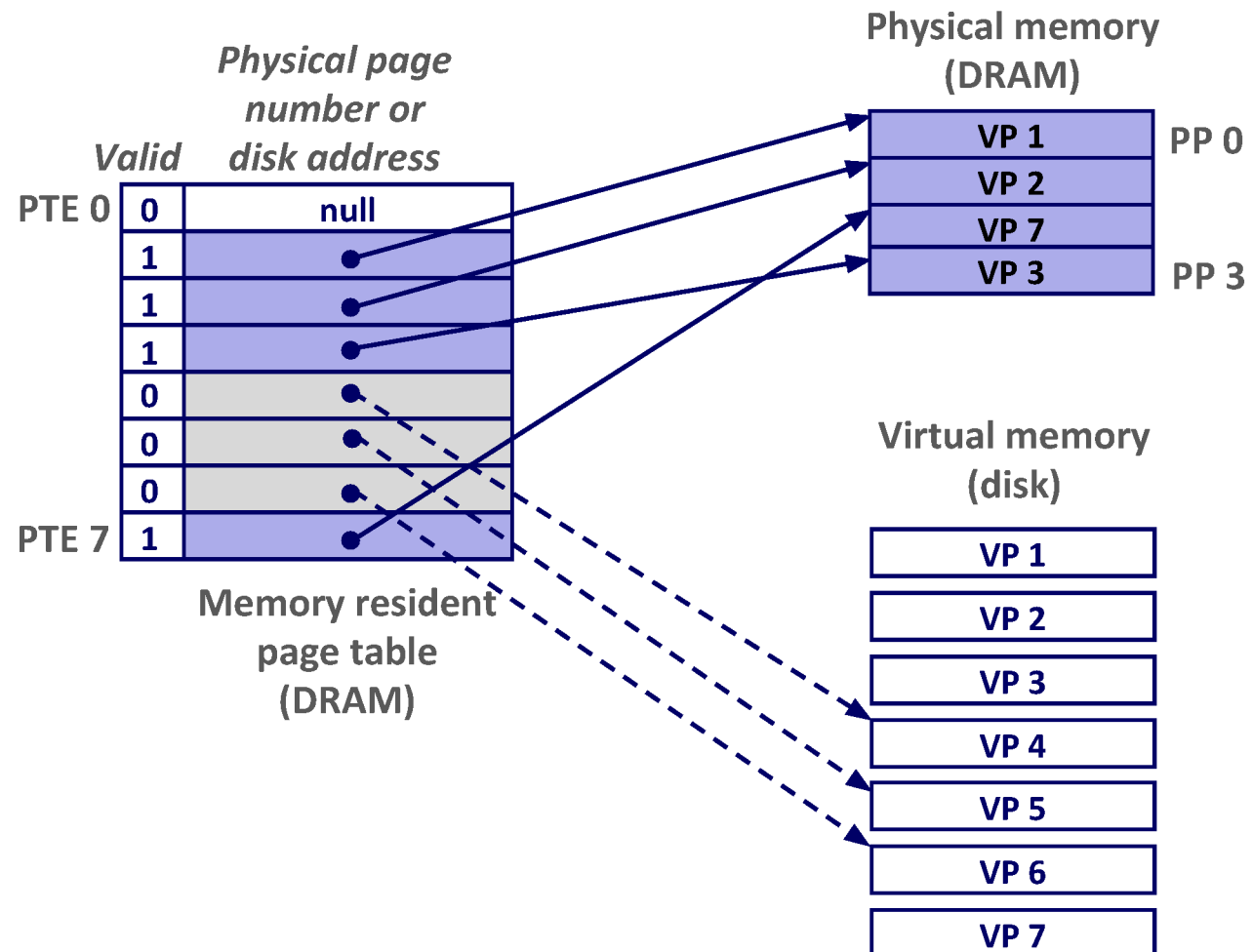
- Page miss causes page fault (an exception)
- Page fault handler selects a victim to be evicted (here VP 4)
- Offending instruction is restarted: page hit!



Key point: Waiting until the miss to copy the page to DRAM is known as *demand paging*

Allocating Pages

- Allocating a new page (VP 5) of virtual memory.



Locality to the Rescue Again!

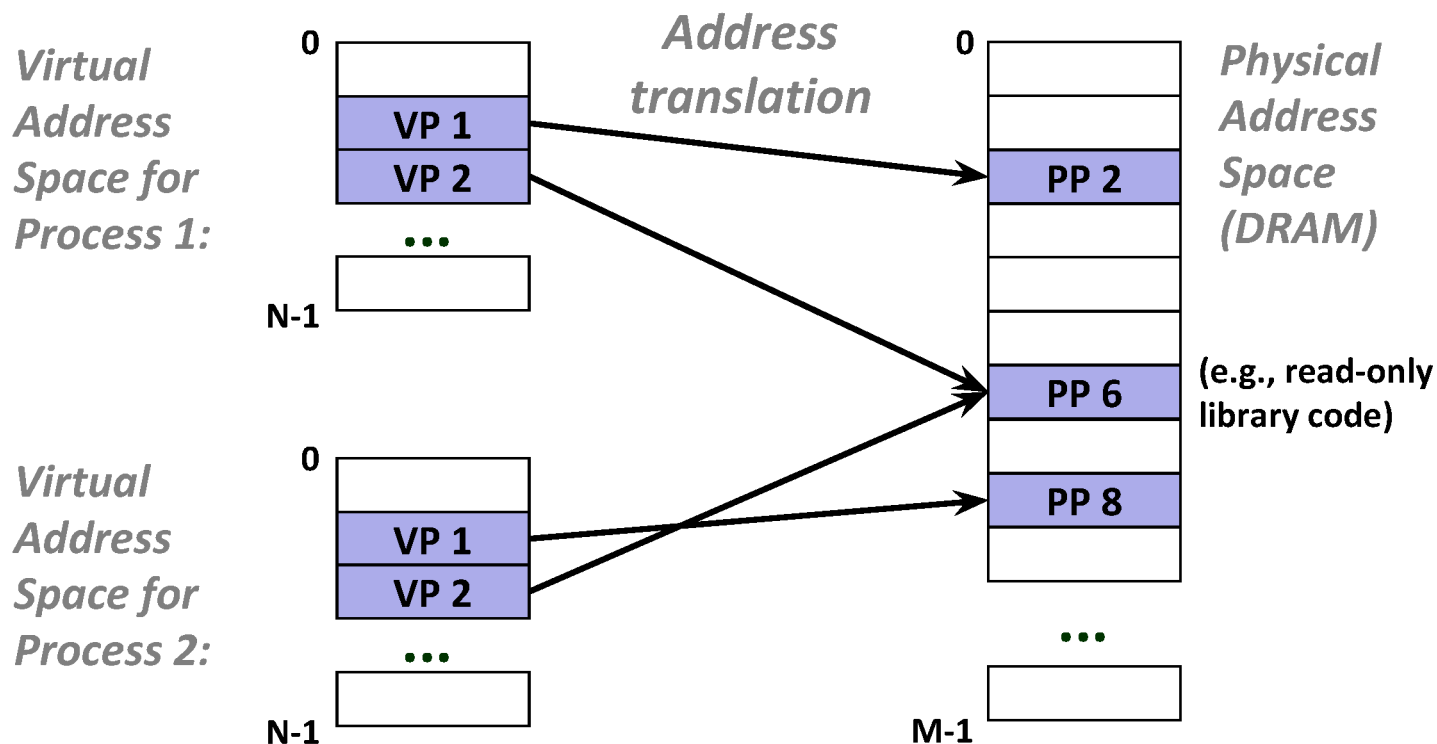
- Virtual memory seems terribly inefficient, but it works because of locality.
- At any point in time, programs tend to access a set of active virtual pages called the *working set*
 - Programs with better temporal locality will have smaller working sets
- If (working set size < main memory size)
 - Good performance for one process after compulsory misses
- If (SUM(working set sizes) > main memory size)
 - *Thrashing*: Performance meltdown where pages are swapped (copied) in and out continuously

Today

- Address spaces
- VM as a tool for caching
- **VM as a tool for memory management**
- VM as a tool for memory protection
- Address translation

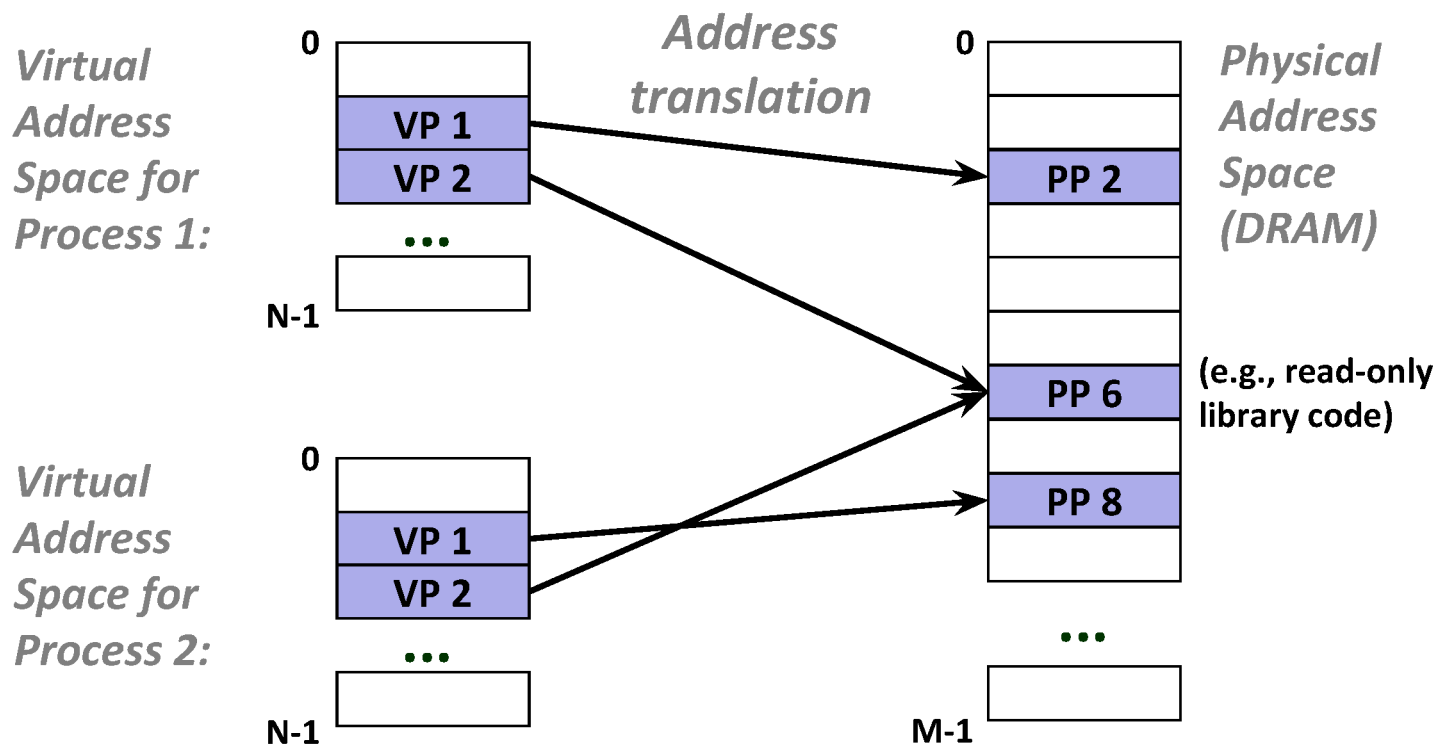
VM as a Tool for Memory Management

- **Key idea: each process has its own virtual address space**
 - It can view memory as a simple linear array
 - Mapping function scatters addresses through physical memory
 - Well-chosen mappings can improve locality



VM as a Tool for Memory Management

- **Simplifying memory allocation**
 - Each virtual page can be mapped to any physical page
 - A virtual page can be stored in different physical pages at different times
- **Sharing code and data among processes**
 - Map virtual pages to the same physical page (here: PP 6)



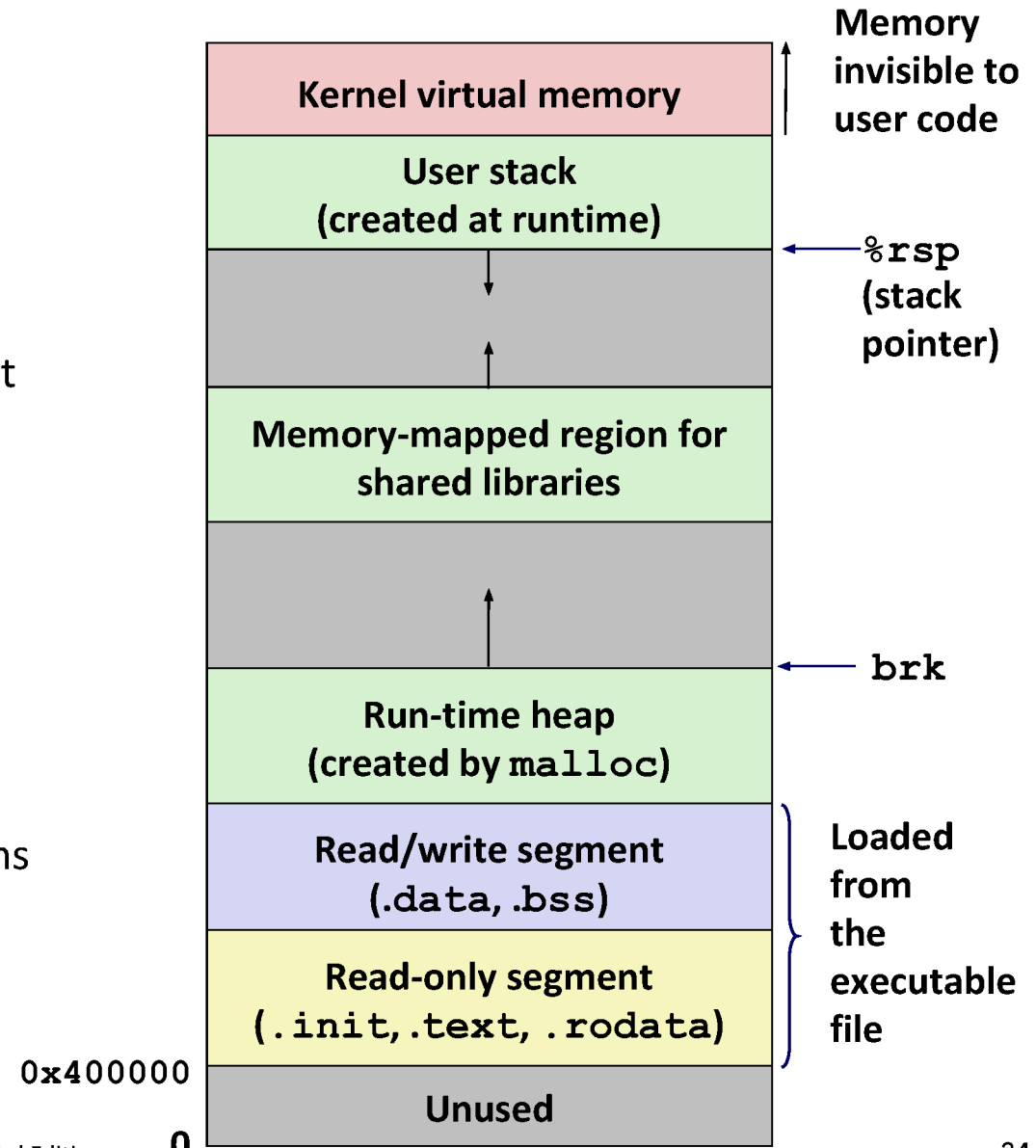
Simplifying Linking and Loading

■ Linking

- Each program has similar virtual address space
- Code, data, and heap always start at the same addresses.

■ Loading

- `execve` allocates virtual pages for `.text` and `.data` sections & creates PTEs marked as invalid
- The `.text` and `.data` sections are copied, page by page, on demand by the virtual memory system

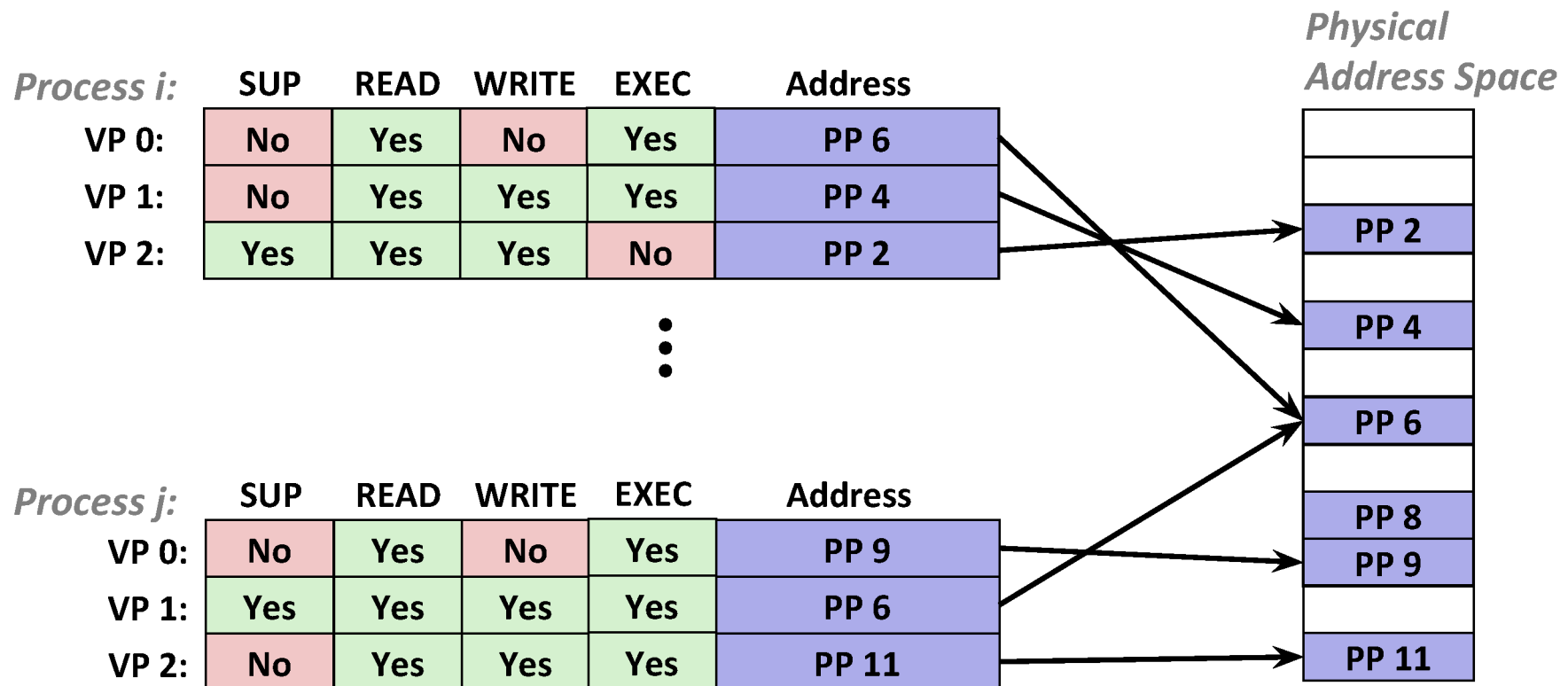


Today

- Address spaces
- VM as a tool for caching
- VM as a tool for memory management
- **VM as a tool for memory protection**
- Address translation

VM as a Tool for Memory Protection

- Extend PTEs with permission bits
- MMU checks these bits on each access



Today

- Address spaces
- VM as a tool for caching
- VM as a tool for memory management
- VM as a tool for memory protection
- **Address translation**

VM Address Translation

■ Virtual Address Space

- $V = \{0, 1, \dots, N-1\}$

■ Physical Address Space

- $P = \{0, 1, \dots, M-1\}$

■ Address Translation

- $MAP: V \rightarrow P \cup \{\emptyset\}$

- For virtual address a :

- $MAP(a) = a'$ if data at virtual address a is at physical address a' in P

- $MAP(a) = \emptyset$ if data at virtual address a is not in physical memory

- Either invalid or stored on disk

Summary of Address Translation Symbols

■ Basic Parameters

- $N = 2^n$: Number of addresses in virtual address space
- $M = 2^m$: Number of addresses in physical address space
- $P = 2^p$: Page size (bytes)

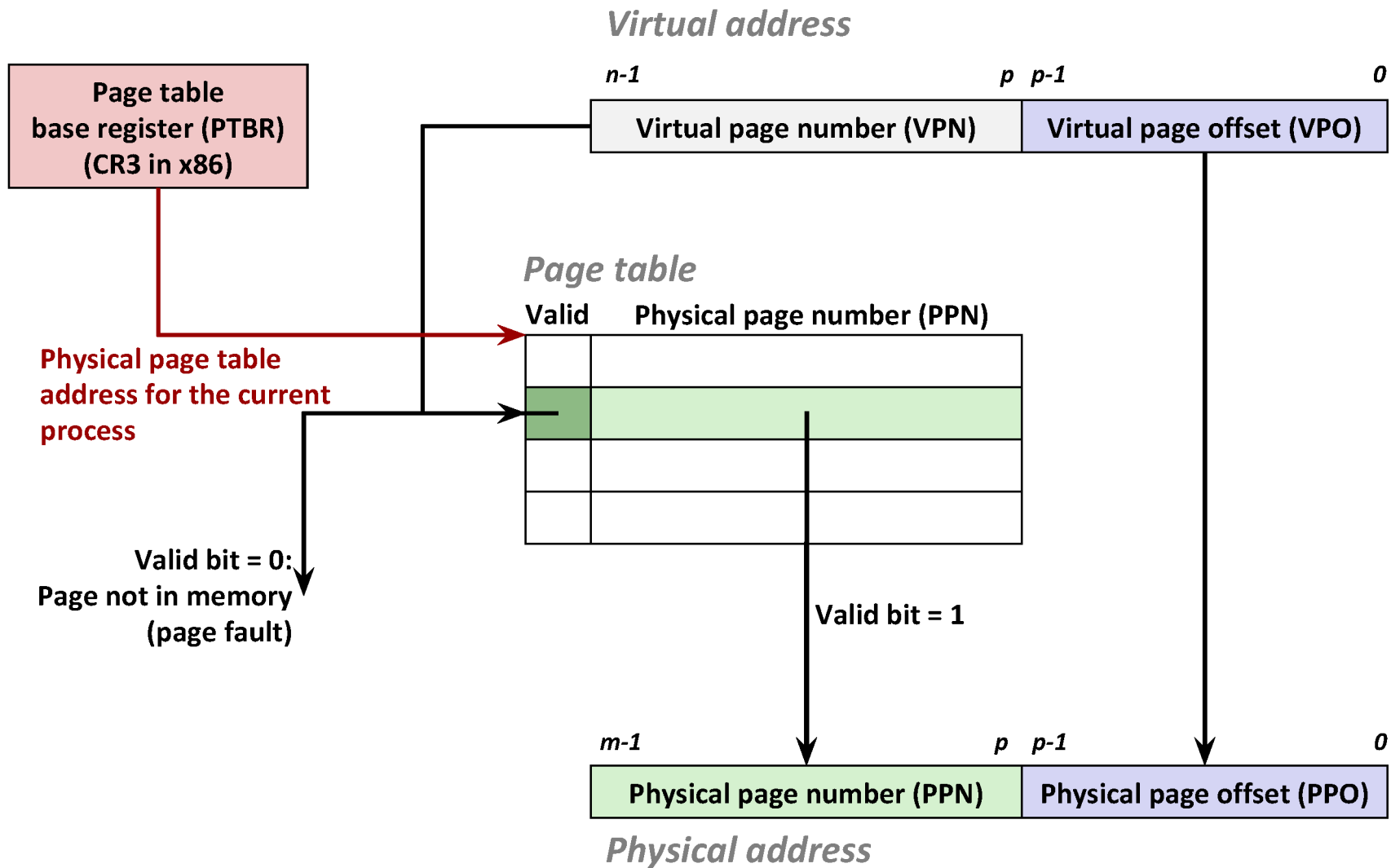
■ Components of the virtual address (VA)

- TLBI: TLB index
- TLBT: TLB tag
- VPO: Virtual page offset
- VPN: Virtual page number

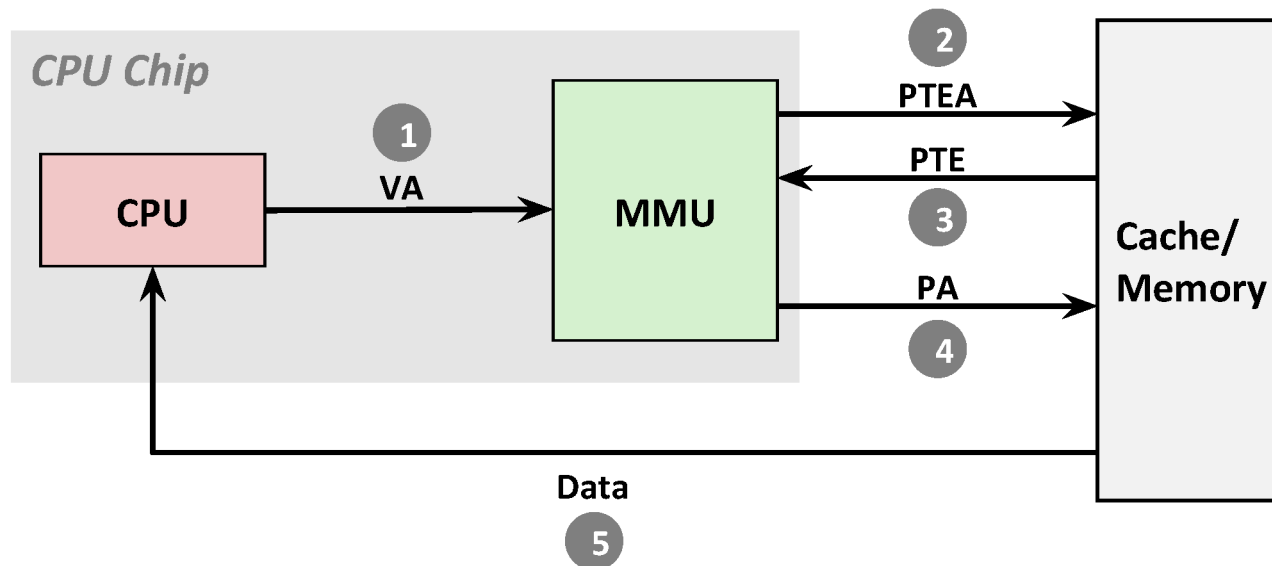
■ Components of the physical address (PA)

- PPO: Physical page offset (same as VPO)
- PPN: Physical page number

Address Translation With a Page Table

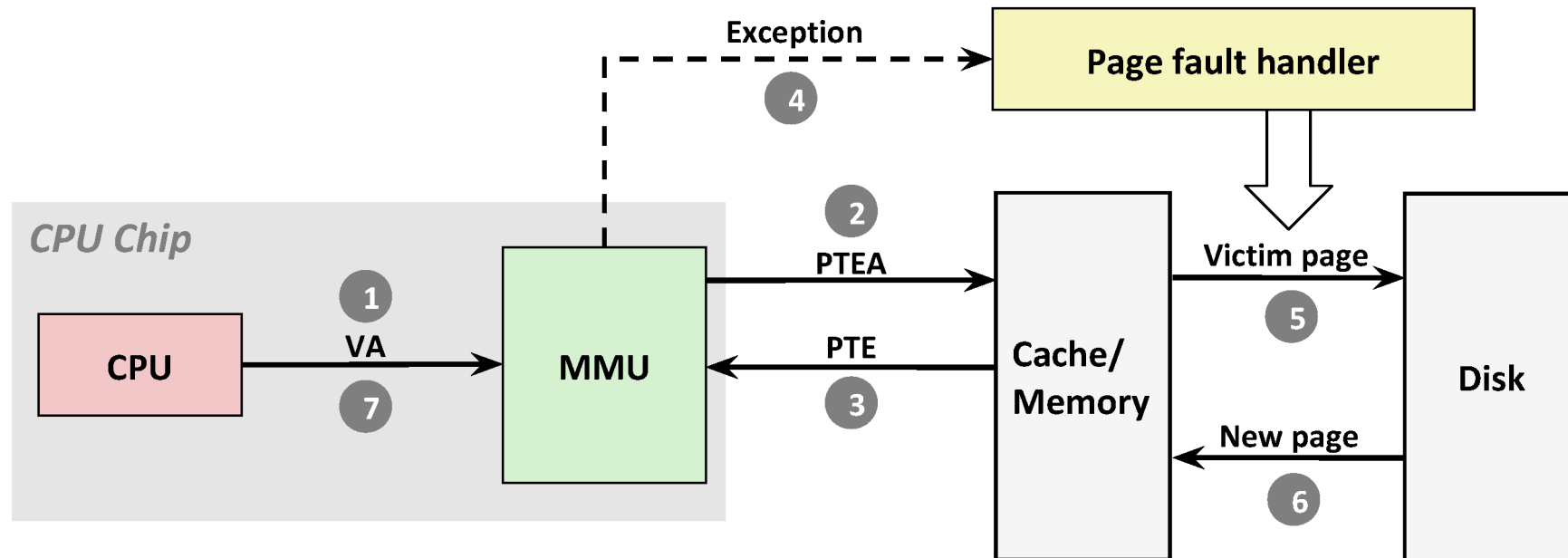


Address Translation: Page Hit



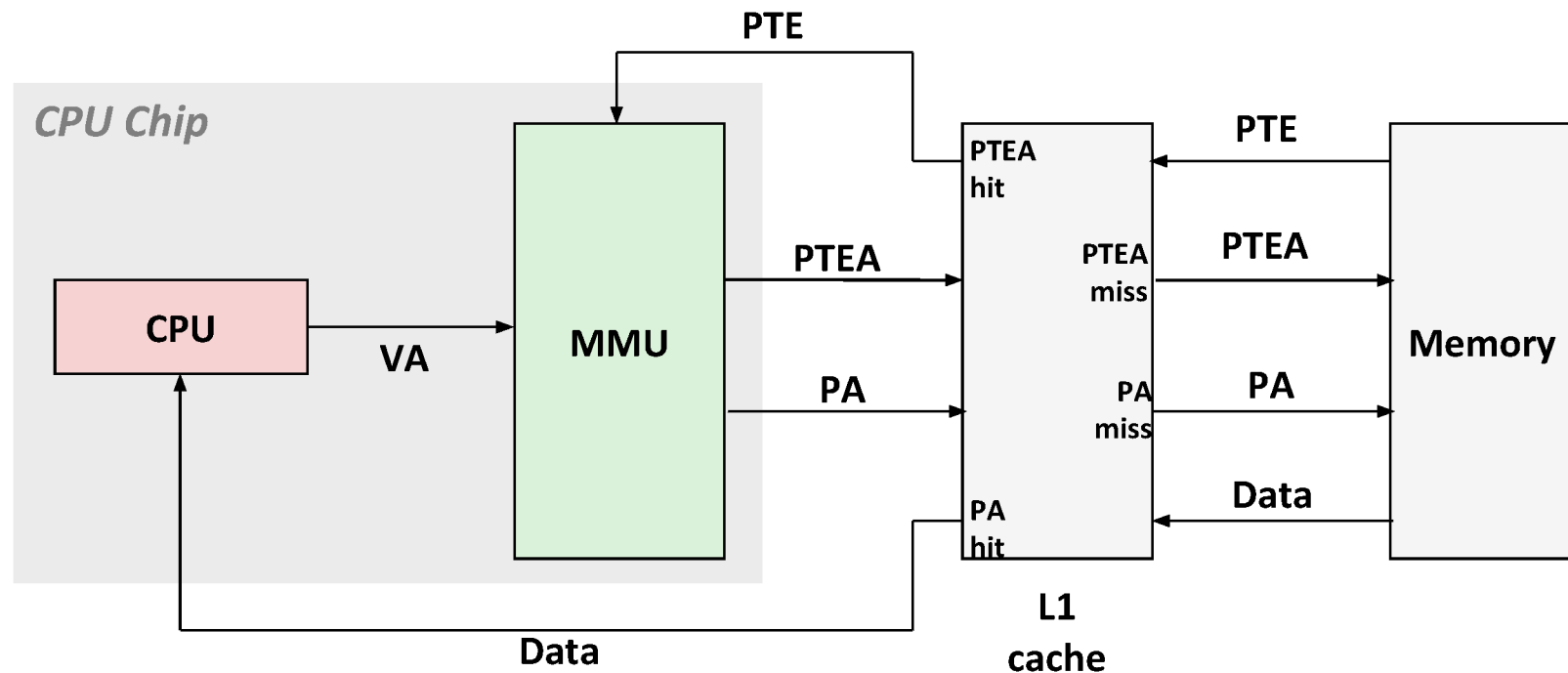
- 1) Processor sends virtual address to MMU
- 2-3) MMU fetches PTE from page table in memory
- 4) MMU sends physical address to cache/memory
- 5) Cache/memory sends data word to processor

Address Translation: Page Fault



- 1) Processor sends virtual address to MMU
- 2-3) MMU fetches PTE from page table in memory
- 4) Valid bit is zero, so MMU triggers page fault exception
- 5) Handler identifies victim (and, if dirty, pages it out to disk)
- 6) Handler pages in new page and updates PTE in memory
- 7) Handler returns to original process, restarting faulting instruction

Integrating VM and Cache



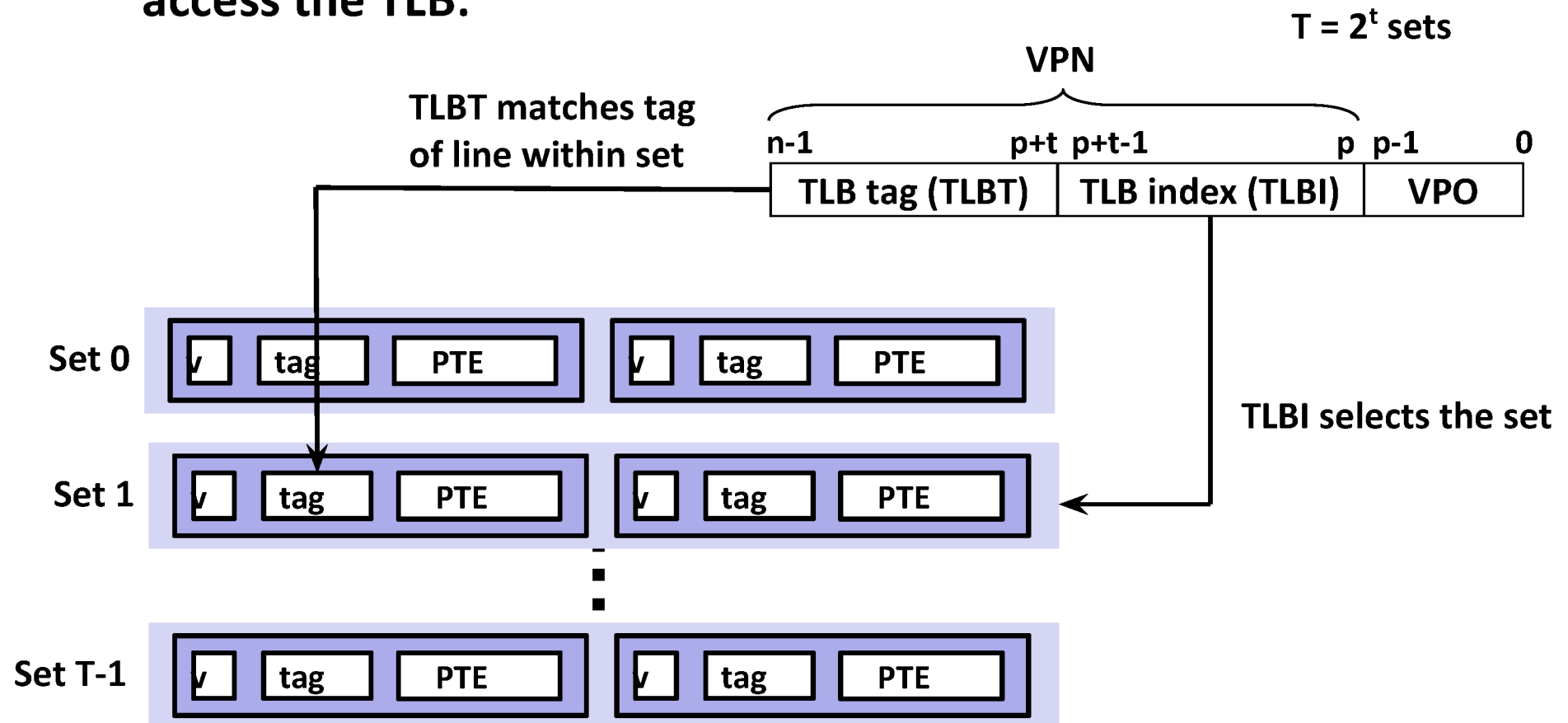
VA: virtual address, PA: physical address, PTE: page table entry, PTEA = PTE address

Speeding up Translation with a TLB

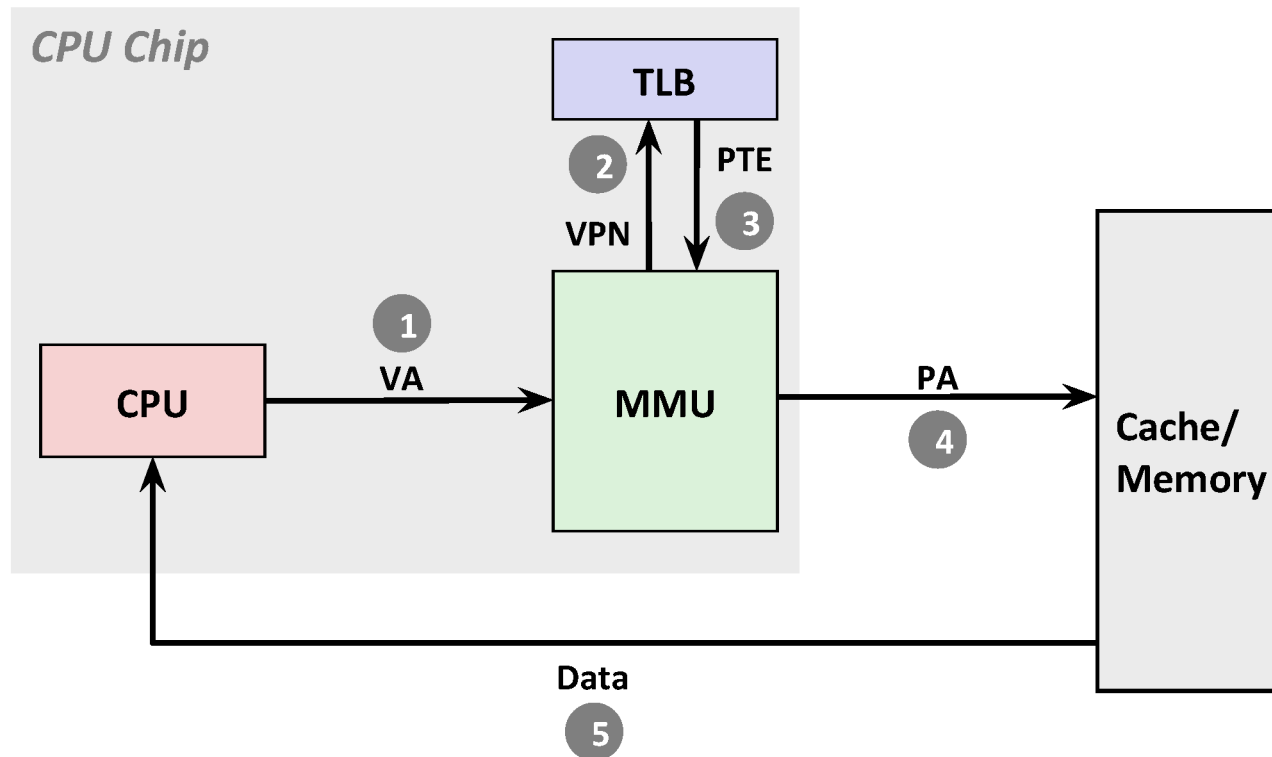
- **Page table entries (PTEs) are cached in L1 like any other memory word**
 - PTEs may be evicted by other data references
 - PTE hit still requires a small L1 delay
- **Solution: *Translation Lookaside Buffer* (TLB)**
 - Small set-associative hardware cache in MMU
 - Maps virtual page numbers to physical page numbers
 - Contains complete page table entries for small number of pages

Accessing the TLB

- MMU uses the VPN portion of the virtual address to access the TLB:

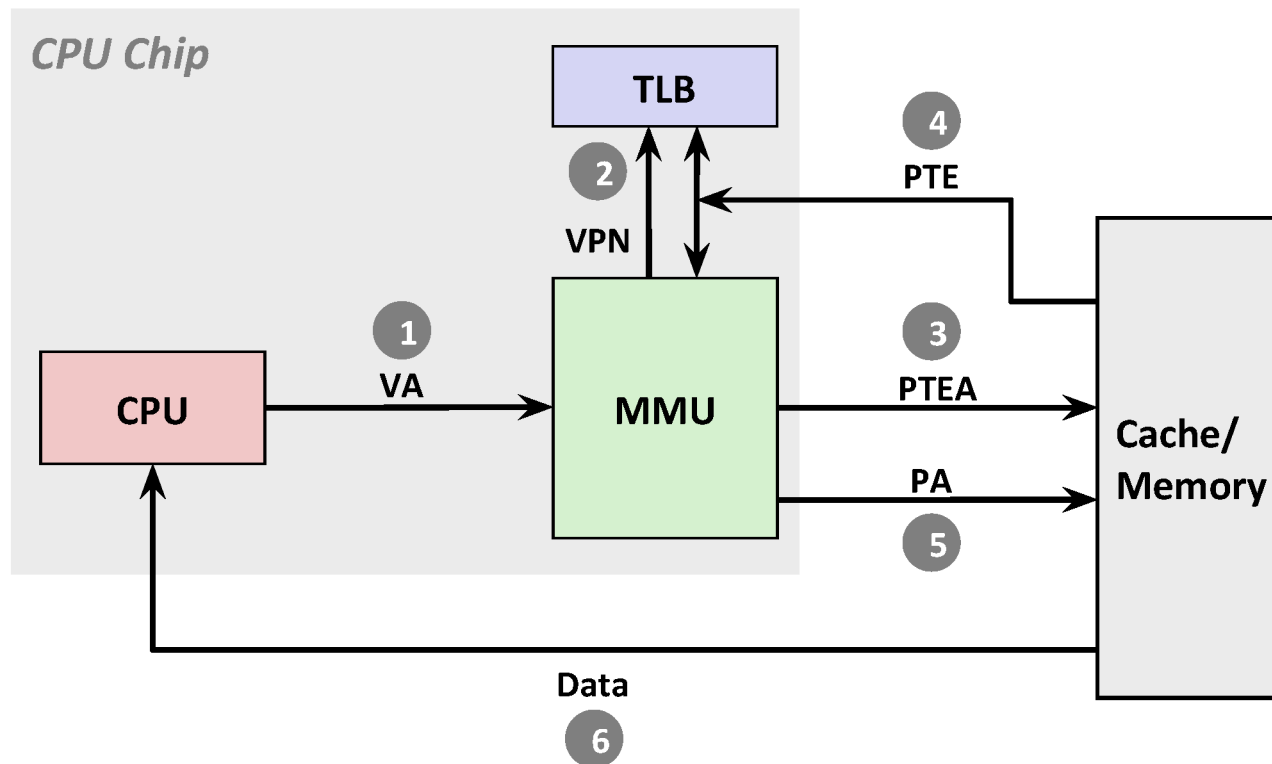


TLB Hit



A TLB hit eliminates a memory access

TLB Miss



A TLB miss incurs an additional memory access (the PTE)

Fortunately, TLB misses are rare. Why?

Multi-Level Page Tables

■ Suppose:

- 4KB (2^{12}) page size, 48-bit address space, 8-byte PTE

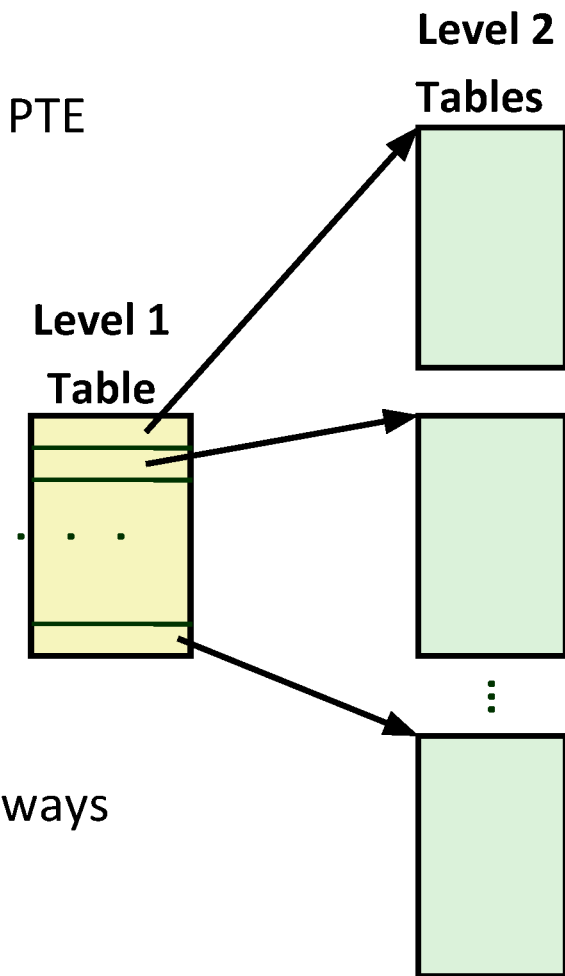
■ Problem:

- Would need a 512 GB page table!
 - $2^{48} * 2^{-12} * 2^3 = 2^{39}$ bytes

■ Common solution: Multi-level page table

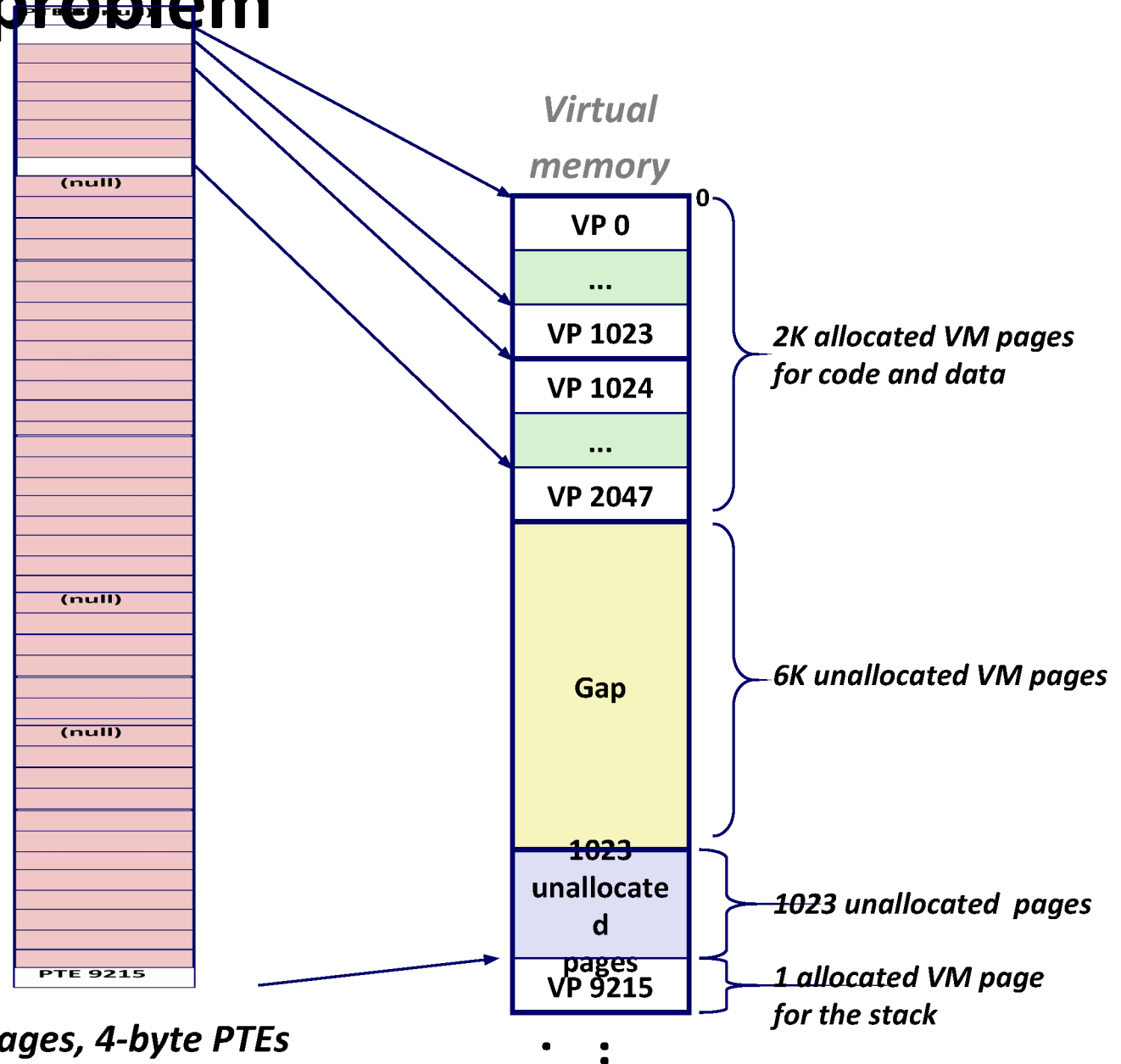
■ Example: 2-level page table

- Level 1 table: each PTE points to a page table (always memory resident)
- Level 2 table: each PTE points to a page (paged in and out like any other data)



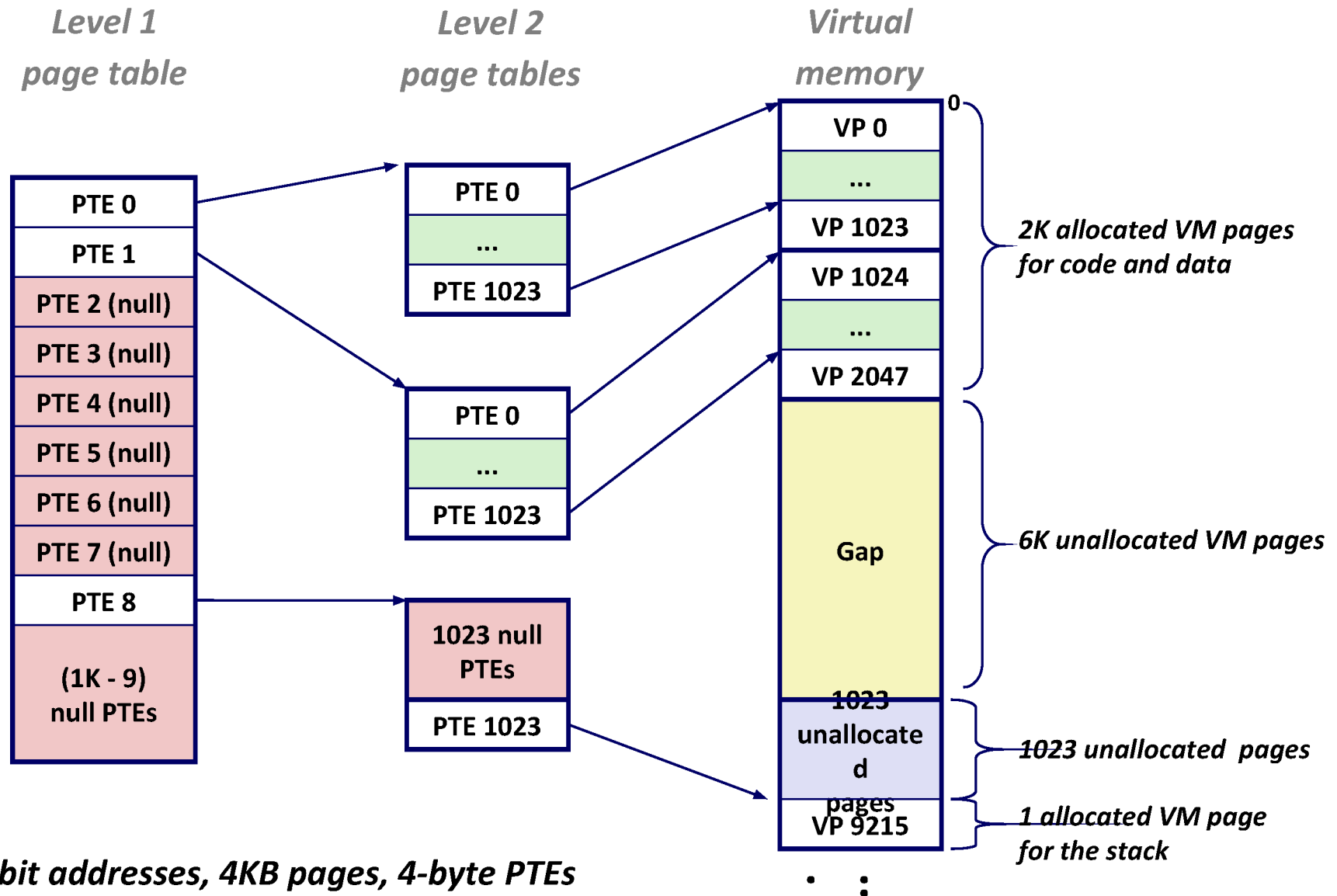
We have a problem

2^{20} Entries of
4 bytes each

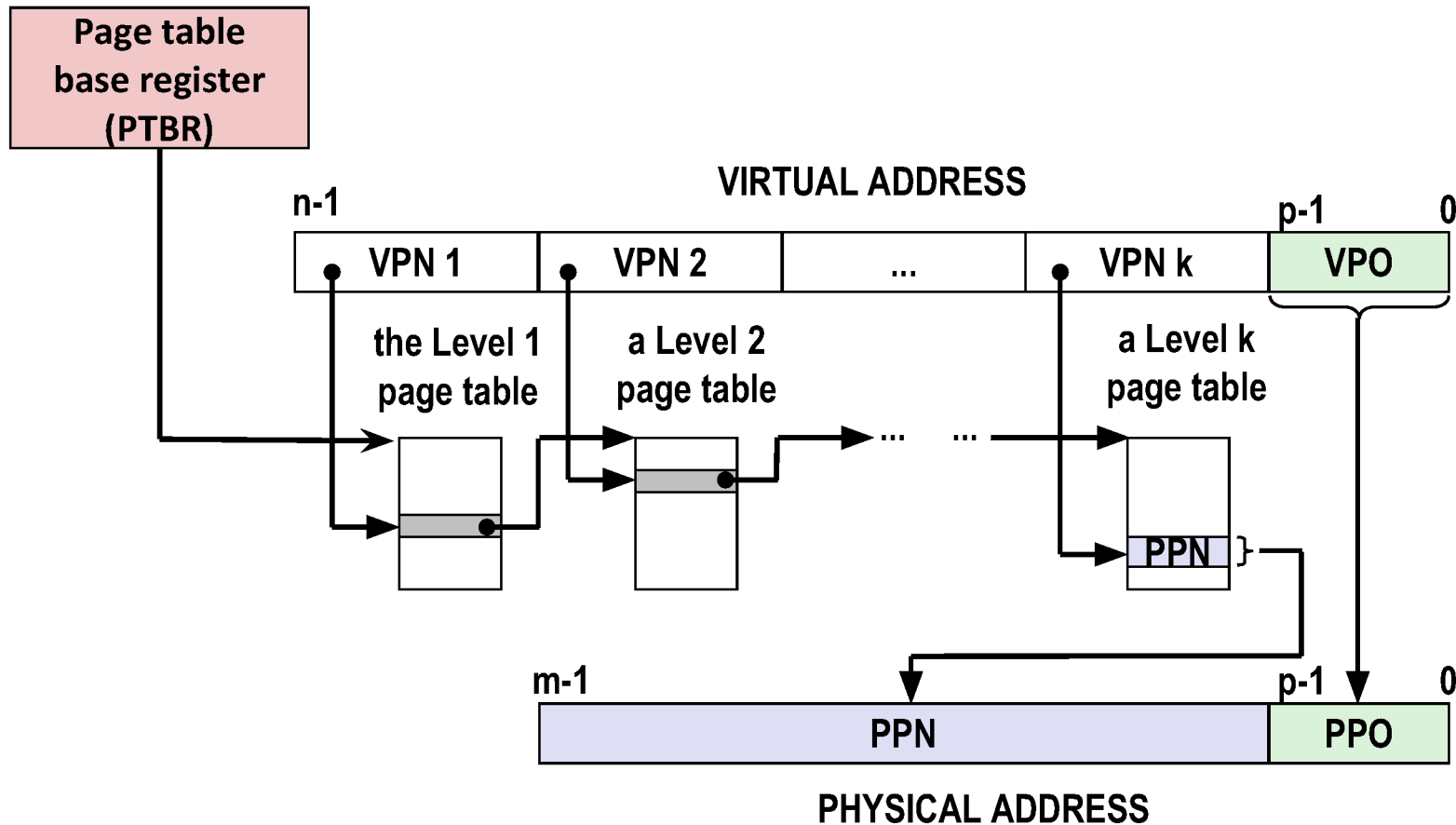


32 bit addresses, 4KB pages, 4-byte PTEs

A Two-Level Page Table Hierarchy



Translating with a k-level Page Table



Summary

- **Programmer's view of virtual memory**
 - Each process has its own private linear address space
 - Cannot be corrupted by other processes

- **System view of virtual memory**
 - Uses memory efficiently by caching virtual memory pages
 - Efficient only because of locality
 - Simplifies memory management and programming
 - Simplifies protection by providing a convenient interpositioning point to check permissions