

## GandALF — Exercise Sheet 9

**Exercise 1.** Let  $\Sigma$  be the set of agents and  $\Gamma \subseteq \Sigma$ . Consider the following ATL formulas:

- (1)  $\langle\langle \Gamma \rangle\rangle Xp \iff \neg \langle\langle \Sigma \setminus \Gamma \rangle\rangle X\neg p$ ,
- (2)  $\langle\langle \Gamma \rangle\rangle F(p \vee q) \iff \langle\langle \Gamma \rangle\rangle F(p) \vee \langle\langle \Gamma \rangle\rangle F(q)$ ,
- (3)  $\langle\langle \Gamma \rangle\rangle F(p \wedge q) \iff \langle\langle \Gamma \rangle\rangle F(p) \wedge \langle\langle \Gamma \rangle\rangle F(q)$ , and
- (4)  $\langle\langle \Gamma \rangle\rangle F(\langle\langle \Gamma \rangle\rangle Fp) \iff \langle\langle \Gamma \rangle\rangle Fp$ .

Which of the above formulas are tautologies? Consider all four semantics: perfect vs. imperfect information, and perfect vs. imperfect recall.

*Remark.* Let  $\mathcal{A}$  be a weighted automaton. We have defined the value of a word  $w$  as the infimum over the values of all accepting runs of  $\mathcal{A}$  on  $w$ . Therefore, if  $\mathcal{A}$  has no accepting runs on  $w$ , the value of  $w$  is the infimum of the empty set, which is  $\infty$ .

**Exercise 2.** Present algorithms deciding the emptiness and the universality problems for weighted automata with the LIMSUP value function. Recall that, if  $\vec{a} = a_1, a_2 \dots$  is a sequence of values from a finite set  $A$ , then  $\text{LIMSUP}(\vec{a})$  is the maximal value from  $A$  that appears infinitely often in  $\vec{a}$ .

**Exercise 3.** Construct a weighted automaton  $\mathcal{A}$  over the SUM value function such that  $\mathcal{A}$  has accepting runs only over words from  $(\#a^*)^*$ , and for all  $w \in (\#a^*)^*$ , which are of the form  $w = \#a^{x_1}\#a^{x_2} \dots \#a^{x_k}\#$  we have:  $\mathcal{L}_{\mathcal{A}}(w) = 0$  if for all  $1 \leq i < k$ , the absolute value of  $x_i - x_{i+1}$  is less or equal to 1, and  $\mathcal{L}_{\mathcal{A}}(w) \leq -1$  otherwise.

**Exercise 4.** Consider a function  $f: \{a, \#\}^\omega \rightarrow \mathbb{R}$  such that  $f(w) < \infty$  if  $w \in (\#a^*)^\omega$ , and  $f(\#a^{x_1}\#a^{x_2} \dots) = \text{LIMAVG}(x_1x_2 \dots)$ . If  $f$  can be expressed by a weighted automaton with the LIMAVG value function, construct such an automaton. Otherwise, show that it is impossible.

**Exercise 5.** Show that  $f_3$  defined for every  $w$  as the edit distance of  $w$  from  $L = a^*b^*$  cannot be computed by a deterministic weighted automaton with the SUM value function.

**Exercise 6.** Present Karp's algorithm computing the minimal mean cycle (prove its correctness and present complexity analysis).